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FILE ID**CDDLIB

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++
TITLE: CDDLIB

CDD Bliss Library

VERSION: 'V04-000'

FACILITY: Common Data Dictionary

ABSTRACT:

This module is the library file used for compiling all other
modules in the CDD facility.

ENVIRONMENT:

AUTHOR: Jeff East, 22-Jan-80

MODIFIED BY:

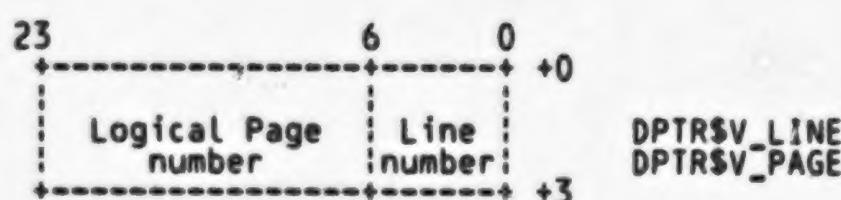
7-May-81 (JAE) Added \$IO_SYNC macro.

-- XTITLE 'CDD Bliss Library'

++ %SBTTL 'STRUCTURE DEFINITIONS'

-- STRUCTURE DEFINITIONS

On-disk Pointer Structure

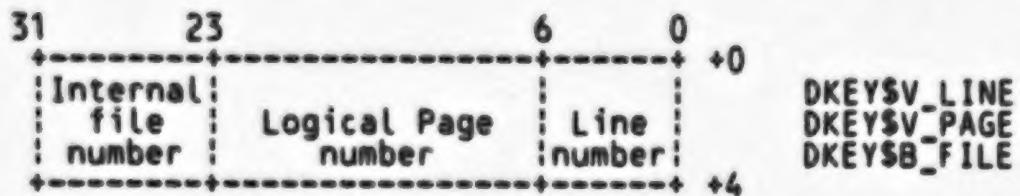


LITERAL
DPTRSS_BLOCK_LENGTH = 3;

MACRO
SDPTR = BLOCK[DPTRSS_BLOCK_LENGTH,BYTE] FIELD (DPTR\$Z_FIELDS)
%:

FIELD DPTR\$Z_FIELDS =
SET
DPTRSV_VALUE = {0, 0, 24, 0},
DPTRSV_LINE = {0, 0, 7, 0},
DPTRSV_PAGE = {0, 7, 17, 0}
TES;

+ In-core Disk Pointer Structure



When a disk pointer is being passed around routines, it is passed as a Disk Key (DKEY), rather than just a disk pointer.

Disk Keys include all the information of a disk pointer, but also include a pointer to their file's FCB.

LITERAL

DKEY\$S_BLOCK_LENGTH = 4;

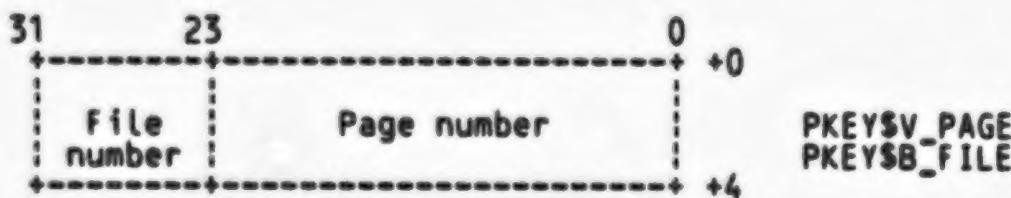
MACRO

SDKEY = BLOCK[DKEY\$S_BLOCK_LENGTH,BYTE] FIELD (DKEY\$Z_FIELDS)
%;

FIELD DKEY\$Z_FIELDS =

SET
DKEY\$V_LINE = [0, 0, 7, 0]
DKEY\$V_PAGE = [0, 7, 17, 0]
DKEY\$B_FILE = [0, 24, 8, 0]
TES;

!+ In-core Page Number Structure



Dictionary page numbers that are passed around the hash table use the Page Key (PKEY) structure to provide both the page and file numbers.

!- LITERAL

```
PKEYSS_BLOCK_LENGTH = 4;
```

MACRO

```
SPKEY = BLOCK[PKEYSS_BLOCK_LENGTH,BYTE] FIELD (PKEYSZ_FIELDS)
```

```
%;
```

FIELD PKEYSZ_FIELDS =

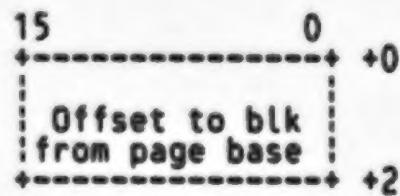
```
SET PKEYSV_PAGE = [0, 0, 24, 0].  
PKEYSB_FILE = [0, 24, 8, 0].
```

```
TES;
```

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Line Index Format

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`LINESW_BLOCK_OFFSET`

Each block on a dictionary page is pointed to by a line index on that page. Disk pointers (DPTR) between the various blocks actually point to the block's line index. The line index is then used to locate the physical dictionary block.

The `PAGES$INDEX` macro allows the programmer to access a line index while it resides on a dictionary page. It also uses the `LINESZ_FIELDS` to access the individual fields within a line index.

5

LITERAL

`LINESS_BLOCK_LENGTH = 2;`

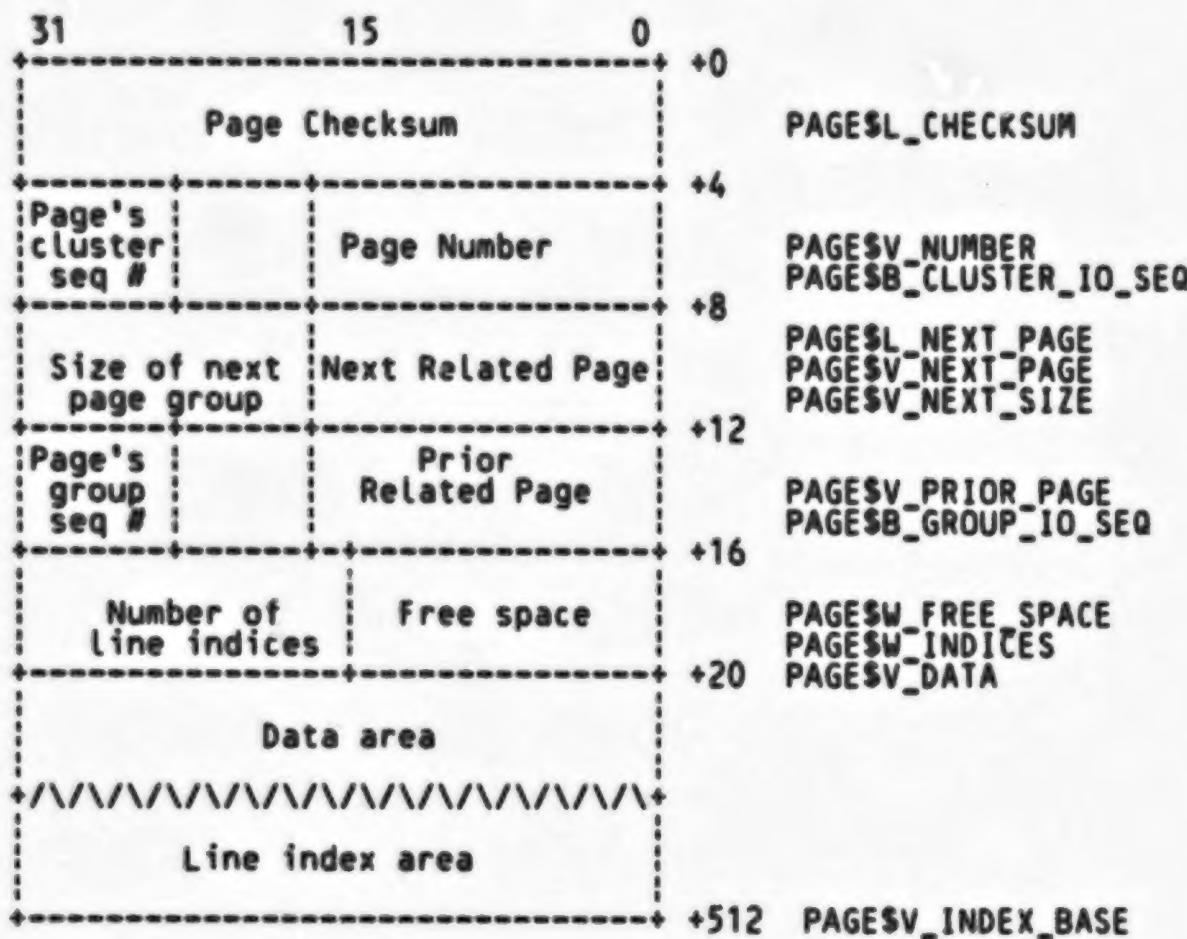
MACRO

`$LINE = BLOCK[LINESS_BLOCK_LENGTH,BYTE] FIELD (LINESZ_FIELDS)`
%;

FIELD LINESZ_FIELDS =

`SET LINESW_BLOCK_OFFSET = [0, 0, 16, 0]`
TES;

Page Format



Each page in the dictionary file starts with a page header. The next and prior related page pointers are used for linking page groups. The following types of page groups exist:

- 1) **Lockable page group.**
Each named entity (and history list) owns exactly one lockable page group. These groups contain all unnamed offspring of the name entry (or history list).

A page group may only be accessed through its portal page. A group's portal page is the page on which the named entity (or history list head) resides. If a group's portal page is locked, then no pages in the group may be accessed.

The **PAGESINDEX** structure is used to access a line index entry on a page.

A page whose checksum is zero is a locked page, and indicates that

the sub-tree below it is incomplete and in a transient state.
Such pages may never be read.

Each page has two page sequence numbers. The cluster sequence number must be the same for all pages in the cluster. The group sequence number must be the same for all pages in an I/O group.

The cluster I/O sequence number is bumped whenever more than one group in the cluster is written. The group I/O sequence number is bumped whenever the group is written. These sequence numbers enable the detection of incomplete cluster unstage operations.

!-

LITERAL

PAGESS_BLOCK_LENGTH = 512;

MACRO

\$PAGE = BLOCK[PAGESS_BLOCK_LENGTH,BYTE] FIELD (PAGESZ_FIELDS)
 %;

FIELD PAGESZ_FIELDS =

SET

PAGESL_CHECKSUM	= [0, 0, 32, 0],
PAGESV_NUMBER	= [4, 0, 17, 0],
PAGESB_CLUSTER_IO_SEQ	= [7, 0, 8, 0],
PAGESL_NEXT_PAGE	= [8, 0, 32, 0],
PAGESV_NEXT_PAGE	= [8, 0, 17, 0],
PAGESV_NEXT_SIZE	= [8, 17, 15, 0],
PAGESV_PRIOR_PAGE	= [12, 0, 17, 0],
PAGESB_GROUP_IO_SEQ	= [15, 0, 8, 0],
PAGESW_FREE_SPACE	= [16, 0, 16, 0],
PAGESW_INDICES	= [18, 0, 16, 0],
PAGESV_DATA	= [20, 0, 0, 0],
PAGESV_INDEX_BASE	= [512, 0, 0, 0]

TES;

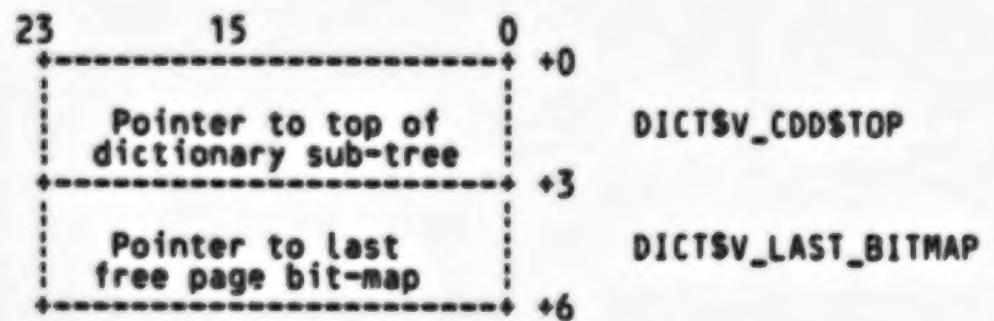
LITERAL

PAGE\$K_BASE = BLOCK[0, PAGESV_DATA;
 PAGESS_BLOCK_LENGTH, BYTE],
PAGE\$K_FREE_SPACE = PAGESS_BLOCK_LENGTH - PAGE\$K_BASE,
PAGESS_INDEX_BASE = BLOCK[0, PAGESV_INDEX_BASE;
 PAGESS_BLOCK_LENGTH, BYTE];

STRUCTURE

PAGESINDEX[i, o, p, s, e] =
(PAGESINDEX+PAGESS_INDEX_BASE+o-(i)*LINESS_BLOCK_LENGTH) <p,s,e>;

Dictionary Header Block



Each dictionary uses page 7 as its dictionary header page.
The dictionary header block immediately follows the page header on page 1.

DICTSV_CDDSTOP points to the root of the tree/sub-tree contained in the dictionary file.

DICTSV_LAST_BITMAP points to the last free page bit-map entry.

LITERAL

```
DICTSS_BLOCK_LENGTH = 6 + PAGESK_BASE;
```

MACRO

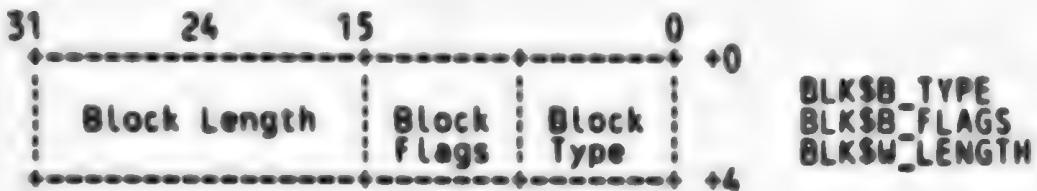
```
$DICT = BLOCK[DICTSS_BLOCK_LENGTH,BYTE] FIELD (PAGESZ_FIELDS,
                                              DICTSZ_FIELDS)
```

%;

FIELD DICTSZ_FIELDS =

```
  SET
    DICTSV_CDDSTOP      = [0+PAGESK_BASE, 0, 24, 0],
    DICTSV_LAST_BITMAP  = [3+PAGESK_BASE, 0, 24, 0]
  TES;
```

Universal Block Header



This block header is used in every non-fixed block on the dictionary pages. Fixed blocks (one which always occur in the same place whenever they exist) do not have this block header.

LITERAL

```
BLKSS_BLOCK_LENGTH = 4;
```

MACRO

```
$BLK = BLOCK[BLKSS_BLOCK_LENGTH,BYTE] FIELD (ACLSZ_FIELDS,
ACLC$Z_FIELDS,
ATT$Z_FIELDS,
BLK$Z_FIELDS,
ELST$Z_FIELDS,
ENT$Z_FIELDS,
FPCBS$Z_FIELDS,
LIST$Z_FIELDS,
LST$Z_FIELDS,
NAME$Z_FIELDS,
NAT$Z_FIELDS,
NNAM$Z_FIELDS,
NODE$Z_FIELDS,
SEG$Z_FIELDS,
SLST$Z_FIELDS,
SSA$Z_FIELDS,
STR$Z_FIELDS,
TEXT$Z_FIELDS)
```

;

```
FIELD BLK$Z_FIELDS =
SET
  BLKSB_TYPE
  BLKSB_FLAGS
  BLK$W_LENGTH
```

```
= [0, 0, 8, 0].
= [1, 0, 8, 0]
= [2, 0, 16, 0]
```

TES;

```
LITERAL
BLK$K_FLAGS
BLK$K_BASE
```

```
= BLOCK[0, BLKSB_TYPE, ,BYTE].
= BLKSS_BLOCK_LENGTH;
```

!

Block types

The ordering and clustering of these block types is significant.

All block types must be contiguous without any holes and bounded by the symbols `BLKSK_TYPE_FIRST` and `BLKSK_TYPE_LAST`.

All node and NAME block types must be contiguous and bounded by the symbols `BLKSK_TYPE_NODE_FIRST` and `BLKSK_TYPE_NODE_LAST`.

To add more node or NAME block types, insert them at the front of the table.

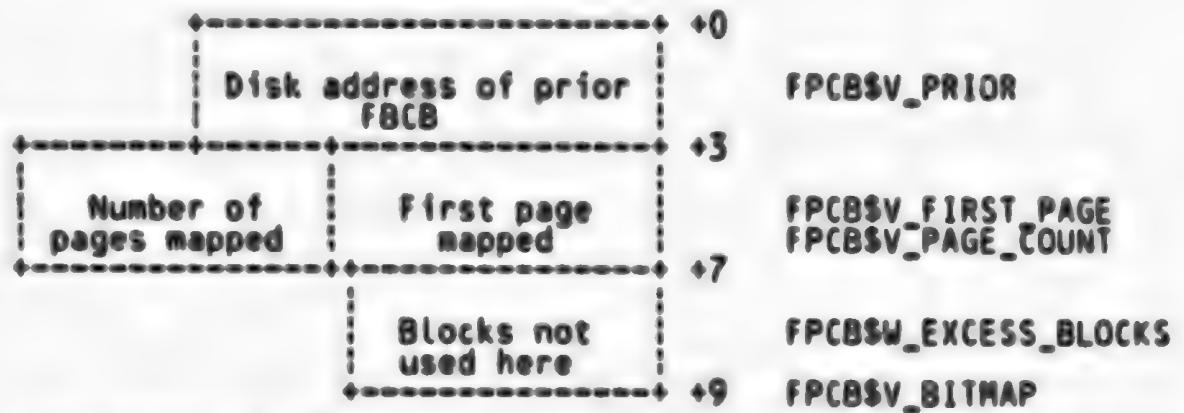
To add any other block type, append them to the end of the table.

CHANGING THE VALUES OF ANY OF THE ACTUAL BLOCK TYPES WILL INVALIDATE EXISTING DICTIONARIES.

LITERAL

<code>BLKSK_TYPE_FIRST</code>	= 101, ! Lowest block type
<code>BLKSK_TYPE_NODE_FIRST</code>	= 101, ! First node or NAM block
<code>BLKSK_TYPE_DIR_NAM</code>	= 101, ! Directory NAM block
<code>BLKSK_TYPE_FIL_NAM</code>	= 102, ! File NAM block
<code>BLKSK_TYPE_TERM_NAM</code>	= 103, ! Terminal NAM block
<code>BLKSK_TYPE_DIR_NODE</code>	= 104, ! Directory node block
<code>BLKSK_TYPE_TERM_NODE</code>	= 105, ! Terminal node block
<code>BLKSK_TYPE_NODE_LAST</code>	= 105, ! Last node or NAM block
<code>BLKSK_TYPE_BITMAP</code>	= 106, ! Free page bitmap
<code>BLKSK_TYPE_ENTITY_ATT</code>	= 107, ! Entity attribute block
<code>BLKSK_TYPE_ENTITY_LIST</code>	= 108, ! Entity list block
<code>BLKSK_TYPE_ENTITY_LIST_ATT</code>	= 109, ! Entity list attribute block
<code>BLKSK_TYPE_NULL_ATT</code>	= 110, ! Null attribute block
<code>BLKSK_TYPE_NUM_ATT</code>	= 111, ! Numeric attribute block
<code>BLKSK_TYPE_SHORT_ATT</code>	= 112, ! Short string attribute block
<code>BLKSK_TYPE_STRING_ATT</code>	= 113, ! String attribute block
<code>BLKSK_TYPE_STRING_LIST</code>	= 114, ! String list block
<code>BLKSK_TYPE_STRING_LIST_ATT</code>	= 115, ! String list attribute block
<code>BLKSK_TYPE_STRING_SEG</code>	= 116, ! String segment block
<code>BLKSK_TYPE_TEXT</code>	= 117, ! Text block
<code>BLKSK_TYPE_ACL</code>	= 118, ! Access control list entry
<code>BLKSK_TYPE_ACCLC</code>	= 119, ! Access control list criterion
<code>BLKSK_TYPE_LAST</code>	= 120, ! Highest block type

Free Page Control Block (FPCB)
 Free Page Bit Map (FPBM)



Free pages in the dictionary file are controlled by the Free Page Control Blocks (FPCB) and the Free Page Bit Map (FPBM). Each free page in the file has its corresponding bit set on.

The disk's cluster size may make it impossible to use all of the allocated blocks. If so, the un-used blocks are tallied in FPCBSW_EXCESS_BLOCKS, and will be used when the file is next extended.

The FPCBs are chained as the dictionary file grows in size.

LITERAL

FPCBSS_BLOCK_LENGTH = BLK\$K_BASE + 9;

STRUCTURE

FPCB[0, P, S, E; BLOCKS] = [FPCBSS_BLOCK_LENGTH+(BLOCKS+7)/8]
 (FPCB+0)<P,S,E>;

FIELD FPCBSZ_FIELDS =

SET

FPCBSV_PRIOR	= [0+BLK\$K_BASE, 0, 24, 0],
FPCBSV_FIRST_PAGE	= [3+BLK\$K_BASE, 0, 17, 0],
FPCBSV_PAGE_COUNT	= [3+BLK\$K_BASE, 17, 15, 0],
FPCBSW_EXCESS_BLOCKS	= [7+BLK\$K_BASE, 0, 16, 0],
FPCBSV_BITMAP	= [9+BLK\$K_BASE, 0, 0, 0]

TES:

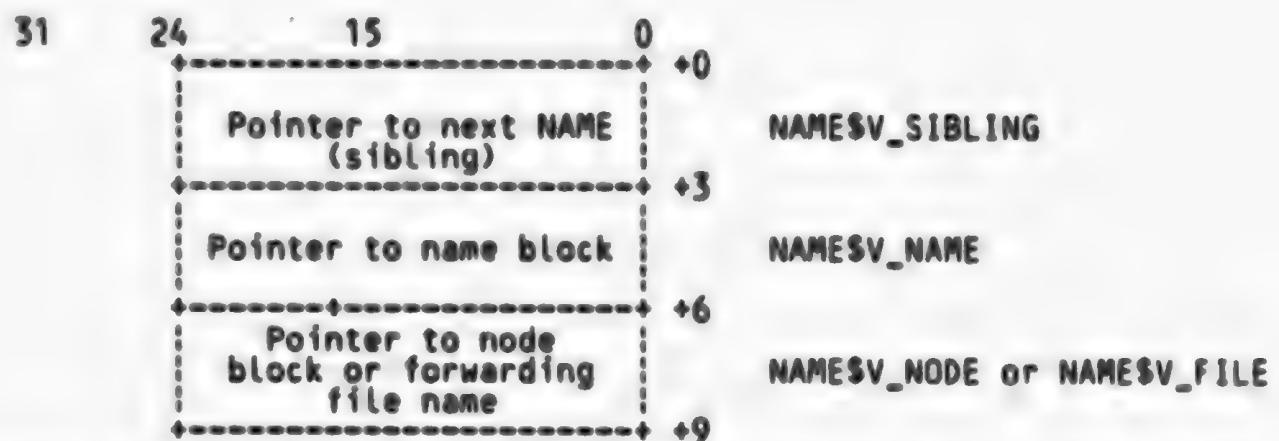
LITERAL

FPCBSK_BASE = FPCBSS_BLOCK_LENGTH,
 FPCBSS_BITMAP = FPCB[0, FPCBSV_BITMAP];

STRUCTURE

FPBM[I] = (FPBM+FPCBSS_BITMAP)<I-1, 1, 0>;

Name Block



There are three types of NAM Blocks, Directory, File, and Terminal.

Directory NAM Block is made up of a Block Header + NAM Block.

File NAM Block is made up of a Block Header + NAM Block.

Terminal NAM Block is made up of a Block Header + NAM Block + Terminal NAM Block.

LITERAL

NAME\$S_BLOCK_LENGTH = 9 + BLK\$K_BASE;

MACRO

SNAME = BLOCK[NAME\$S_BLOCK_LENGTH,BYTE] FIELD (BLK\$Z_FIELDS,
NAME\$Z_FIELDS)

%;

FIELD NAME\$Z_FIELDS =

SET

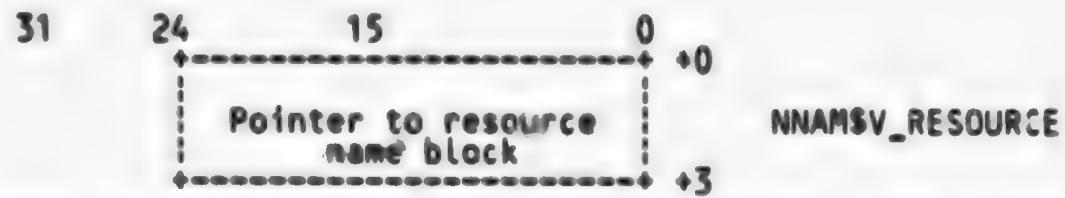
NAME\$V_SIBLING	= [0+BLK\$K_BASE, 0, 24, 0],
NAME\$V_NAME	= [3+BLK\$K_BASE, 0, 24, 0],
NAME\$V_NODE	= [6+BLK\$K_BASE, 0, 24, 0],
NAME\$V_FILE	= [6+BLK\$K_BASE, 0, 24, 0]

TES.

LITERAL

NAME\$K_BASE = NAME\$S_BLOCK_LENGTH;

* Node Nam Block



Node NAM Block is made up of a Block Header + NAM Block + Node NAM Block.

LITERAL

NNAMSS_BLOCK_LENGTH = 3 + NAMESK_BASE;

MACRO

\$NNAM = BLOCK[NNAMSS_BLOCK_LENGTH,BYTE] FIELD (BLK\$Z_FIELDS,
NAMESZ_FIELDS,
NNAMSZ_FIELDS)

%;

FIELD NNAMSZ_FIELDS =

SET

NNAMSV_RESOURCE = [0+NAMESK_BASE, 0, 24, 0]

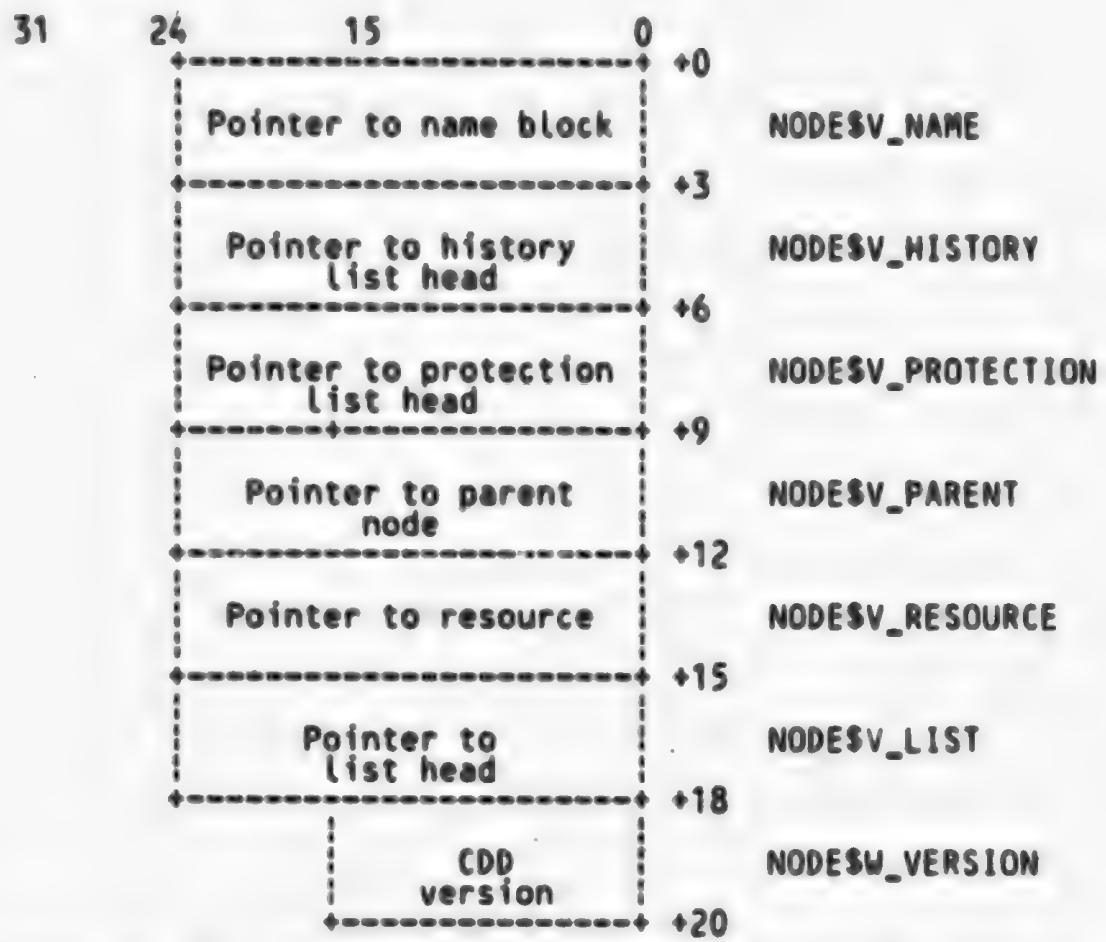
TES;

LITERAL

NNAMSK_BASE

= NNAMSS_BLOCK_LENGTH;

+ Node Block



There are two types of Node Blocks Directory and Terminal.

Directory Node Block is made up of a Block Header + Node Block

Terminal Node Block is made up of a Block Header + Node Block

- LITERAL

```
NODESS_BLOCK_LENGTH = 20+ BLK$K_BASE;
```

- MACRO

```
$NODE = BLOCK[NODESS_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS,  
NODES$Z_FIELDS)
```

```
%:
```

```
FIELD NODES$Z_FIELDS =  
SET
```

```
NODESV_ORDERED = [1, 0, 1, 0]  
NODESV_NAME = [0+BLK$K_BASE, 0, 24, 0];  
NODESV_HISTORY = [3+BLK$K_BASE, 0, 24, 0];
```

NODESV_PROTECTION = [6+BLKSK_BASE, 0, 24, 0],
NODESV_PARENT = [9+BLKSK_BASE, 0, 24, 0],
NODESV_RESOURCE = [12+BLKSK_BASE, 0, 24, 0],
NODESV_LIST = [15+BLKSK_BASE, 0, 24, 0],
NODESV_VERSION = [18+BLKSK_BASE, 0, 16, 0]

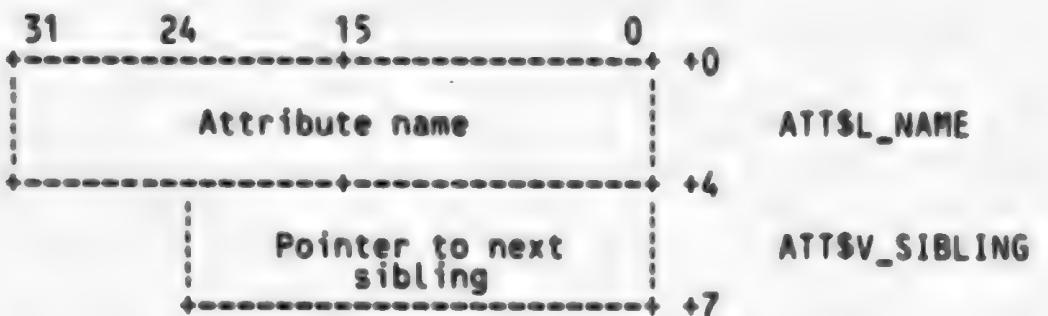
TES:

LITERAL
NODESK_BASE = NODESS_BLOCK_LENGTH;

The following flag occurs in the BLKSB_FLAGS field of
a directory node.

LITERAL
NODESM_ORDERED = 1^1 - 1^0: ! List elements are sorted by name

!+ Common Attribute Block



Each attribute has the common attribute block immediately following the universal block header.

There are six types of Attribute Blocks Entity Attribute Block, List Attribute Block, Numeric Attribute Block, Null Attribute Block, Short String Attribute Block, and String Attribute Block.

Entity Attribute Block is made up of a Block Header + Attribute Block + Entity Attribute Block.

List Attribute Block is made up of a Block Header + Attribute Block + List Attribute Block.

Numeric Attribute Block is made up of a Block Header + Attribute Block + Numeric Attribute Block.

Null Attribute Block is made up of a Block Header + Attribute Block.

Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

!-

LITERAL

ATTSS_BLOCK_LENGTH = 7 + BLKSZ_BASE;

MACRO

SATT = BLOCK[ATTSS_BLOCK_LENGTH,BYTE] FIELD (ATTsz_FIELDS,
BLKSz_FIELDS)
%;FIELD ATTsz_FIELDS =
SETATTSL_NAME = [0+BLKSZ_BASE, 0, 32, 0].
ATTsv_SIBLING = [4+BLKSZ_BASE, 0, 24, 0].

TES;

LITERAL

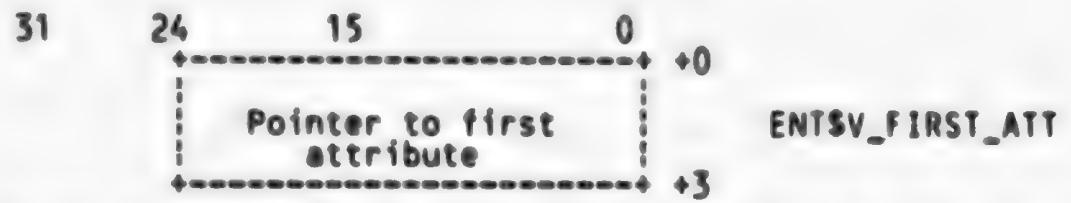
CDDL1B.B32;1

16-SEP-1984 16:58:51.80 Page 17

ATT\$K_BASE

= ATT\$S_BLOCK_LENGTH;

Entity Attribute Block



Entity Attribute Block is made up of a Block Header + Attribute Block
+ Entity Attribute Block.

LITERAL
ENT\$S_BLOCK_LENGTH = 3 + ATT\$K_BASE;

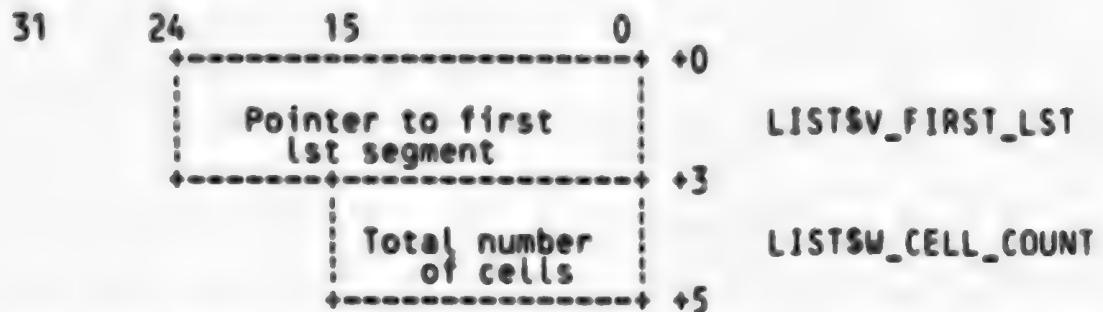
MACRO
SENT = BLOCK[ENT\$S_BLOCK_LENGTH,BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
ENT\$Z_FIELDS)

%;

FIELD ATT\$Z_FIELDS =
SET ENT\$V_FIRST_ATT = [0+ATT\$K_BASE, 0, 24, 0]
TES;

LITERAL
ATT\$K_BASE = ENT\$S_BLOCK_LENGTH;

List Attribute Block



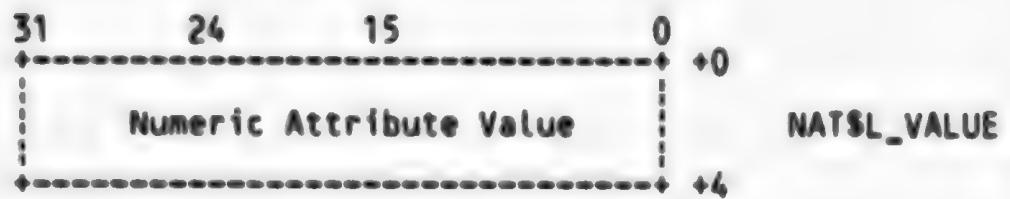
- List Attribute Block is made up of a Block Header + Attribute Block

LITERAL
LISTSS_BLOCK_LENGTH = 5 + ATTSK_BASE;

FIELD LIST\$Z_FIELDS =
SET
 LIST\$V_FIRST_LST = [0+ATTR\$K_BASE, 0, 24, 0],
 LIST\$W_CELL_COUNT = [3+ATTR\$K_BASE, 0, 16, 0]
TES:

LITERAL LIST\$K_BASE = LIST\$S_BLOCK_LENGTH;

Numeric Attribute Block



Numeric Attribute Block is made up of a Block Header + Attribute Block
+ Numeric Attribute Block.

Numeric attribute block contains the
value of the numeric attribute.

LITERAL
NATSS_BLOCK_LENGTH = 4 + ATT\$K_BASE;

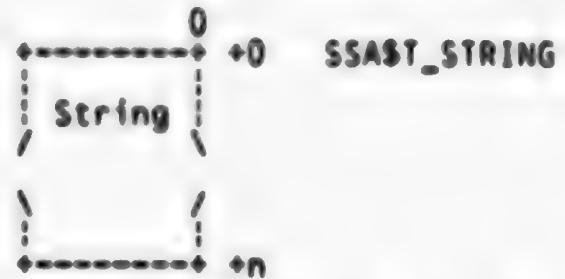
MACRO
SNAT = BLOCK[NATSS_BLOCK_LENGTH,BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
NAT\$Z_FIELDS)
%;

FIELD NAT\$Z_FIELDS =
SET
NATSL_VALUE = [0+ATT\$K_BASE,0,32,0]
TES;

LITERAL
NAT\$K_BASE = NATSS_BLOCK_LENGTH;

Short String Attribute Block

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Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

Strings whose total length is between 0 and 255 bytes are usually stored using the short string attribute block. If a string is too long to be stored in an SSA, then it is stored using a normal string attribute block (STR).

LITERAL

```
SSASS_BLOCK_LENGTH = 0 + ATTRK_BASE;
```

MACRO

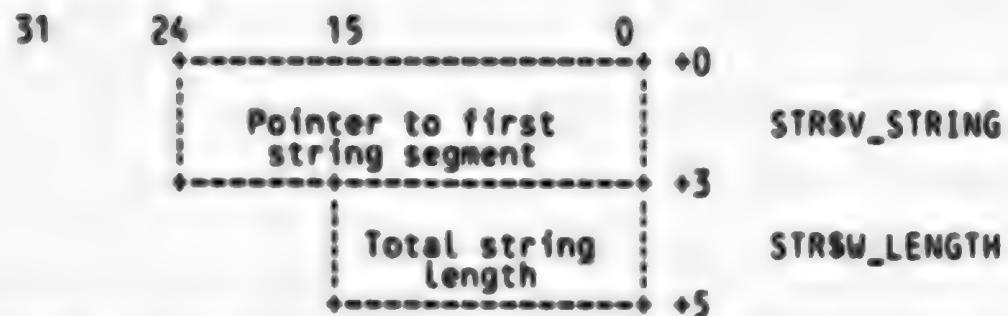
```
SSSA = BLOCK[SSASS_BLOCK_LENGTH, BYTE] FIELD (ATTRS_FIELDS,  
BLKSZ_FIELDS,  
SSASZ_FIELDS)
```

%

```
FIELD SSASZ_FIELDS =  
SET  
  SSABT_STRING      = [0+ATTRK_BASE, 0, 0, 0]  
TES:
```

```
LITERAL  
  SSASK_BASE      = SSASS_BLOCK_LENGTH;
```

String Attribute Block



String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

The string is broken into one or more segments. The string segments are stored in SEG blocks.

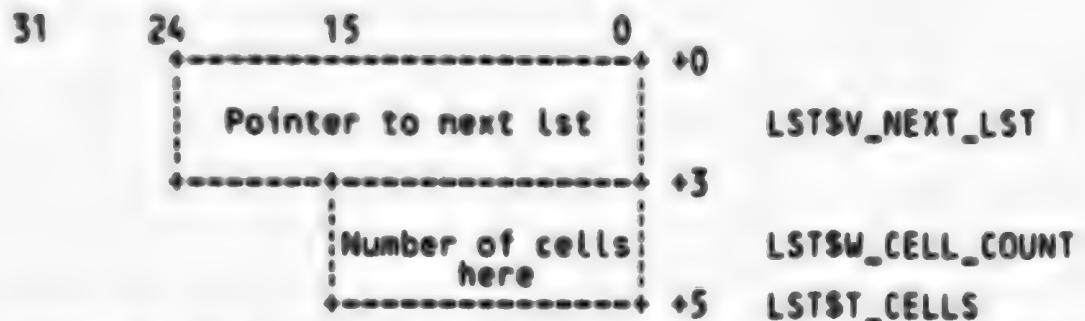
LITERAL
STRSS_BLOCK_LENGTH = 5 * ATTSK_BASE;

MACRO **8STR** = **BLOCK[STRSS_BLOCK_LENGTH,BYTE]** **FIELD** {**ATTSZ_FIELDS,**
BLKSZ_FIELDS,
STRSZ_FIELDS}

%:
FIELD STRSZ_FIELDS =
SET
 STRSV_STRING
 STRSW_LENGTH
TES:
= [0+ATTBK_BASE, 0, 24, 0].
= [3+ATTBK_BASE, 0, 16, 0].

LITERAL
STR8K_BASE = STR8S_BLOCK_LENGTH;

* LST Segment Block



There are two types of LST Segment Blocks Entity List Block and String List Attribute Block.

Entity List Block is made up of a Block Header + LST Block + Entity List Block.

String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.

LITERAL

```
LS15S_BLOCK_LENGTH = 5 + BLKSK_BASE;
```

MACRO

```
LS15 = BLOCK[LS15S_BLOCK_LENGTH,BYTE] FIELD (BLKSK_FIELDS,
                                             LS15Z_FIELDS)
```

%

FIELD LS15Z_FIELDS =

SET

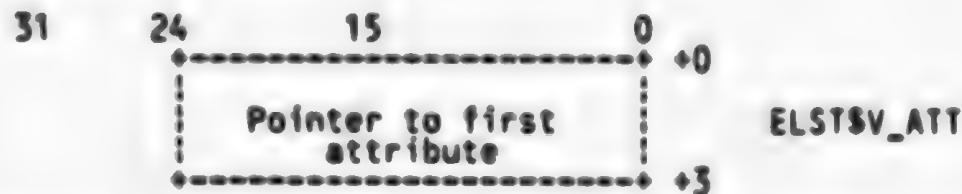
LS15V_NEXT_LST	= [0 + BLKSK_BASE, 0, 24, 0],
LS15W_CELL_COUNT	= [3 + BLKSK_BASE, 0, 16, 0],
LS15T_CELLS	= [5 + BLKSK_BASE, 0, 0, 0]

TES:

LITERAL

```
BLKSK_BASE = LS15S_BLOCK_LENGTH;
```

Entity List Block



Entity List Block is made up of a Block Header + LST Block + Entity List Block.

LITERAL
ELSTSS_BLOCK_LENGTH = 3;

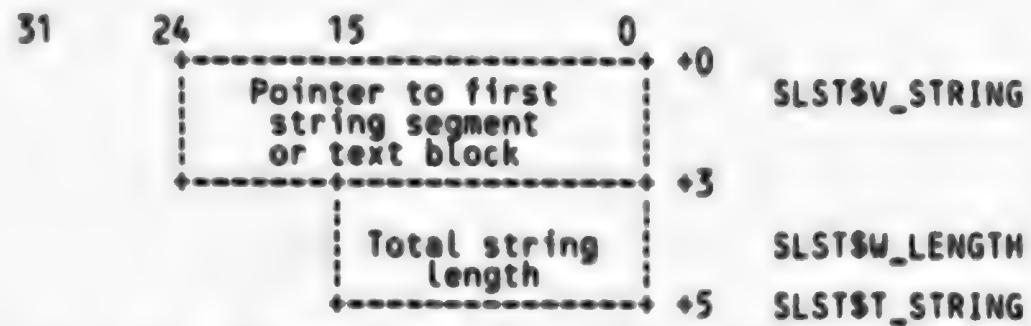
MACRO
SELST = BLOCK[ELSTSS_BLOCK_LENGTH,BYTE] FIELD (ELSTSZ_FIELDS)
%;

FIELD ELSTSZ_FIELDS =
SET
ELSTSV_ATT = [0, 0, 24, 0]
TES;

LITERAL
ELSTSK_BASE = ELSTSS_BLOCK_LENGTH;

STRUCTURE
ELSTSELM[I] =
(ELSTSELM + LSTSK_BASE + (I * ELSTSS_BLOCK_LENGTH));

* String List Attribute Block



String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.

LITERAL
SLSTSS_BLOCK_LENGTH = 5;

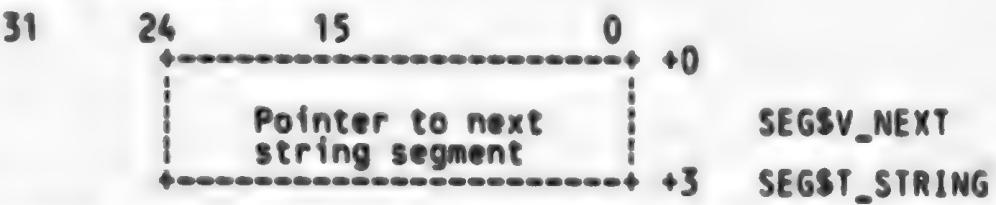
MACRO
SSLST = BLOCK[SLSTSS_BLOCK_LENGTH,BYTE] FIELD (SLSTSZ_FIELDS)
%;

FIELD SLSTSZ_FIELDS =
SET
SLSTSV_STRING = [0, 0, 24, 0],
SLSTSW_LENGTH = [3, 0, 16, 0],
SLSTST_STRING = [5, 0, 0, 0]
TES;

LITERAL
SLSTSK_BASE = SLSTSS_BLOCK_LENGTH;

STRUCTURE
SLSTSELM[] =
(SLSTSELM + LSTSK_BASE + (I * SLSTSS_BLOCK_LENGTH));

String Segment Block



String Segment Block is made up of a Block Header + String Segment Block.

Each string attribute points to zero or more string segment blocks. Each block contains a portion of the whole string. The final string is found by concatenating all the string segments together.

LITERAL
SEGSS_BLOCK_LENGTH = 3 + BLKSK_BASE;

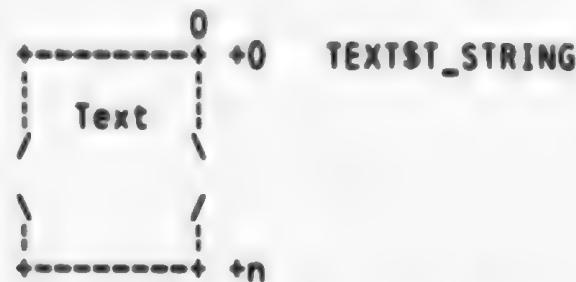
MACRO \$SEG = BLOCK[SEGSS_BLOCK_LENGTH,BYTE] FIELD (BLKSZ_FIELDS,
SEGSS_FIELDS)
%;

FIELD SEG\$Z_FIELDS =
SET SEG\$V_NEXT, SEG\$T_STRING
TES: = [0+BLK\$K_BASE, 0, 24, 0],
= [3+BLK\$K_BASE, 0, 0, 0]

LITERAL SEGSK_BASE = SEGSS_BLOCK_LENGTH;

Text Block

31



Text Block is made up of a Block Header + Text Block.

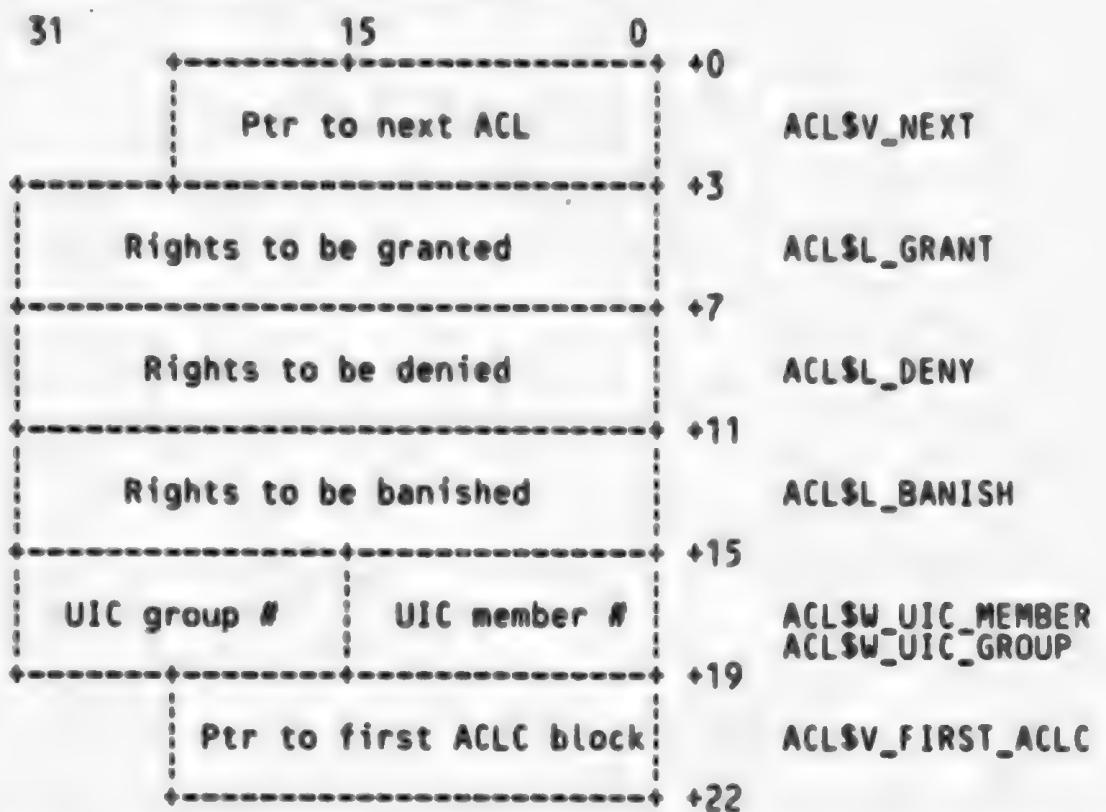
LITERAL
TEXTSS_BLOCK_LENGTH = 0 + BLK\$K_BASE;

MACRO
\$TEXT = BLOCK[TEXTSS_BLOCK_LENGTH,BYTE] FIELD (BLK\$Z_FIELDS,
TEXT\$Z_FIELDS)
%;

FIELD TEXT\$Z_FIELDS =
SET
TEXTST_STRING . = [0+BLK\$K_BASE, 0, 0, 0]
TES;

LITERAL
TEXT\$K_BASE = TEXTSS_BLOCK_LENGTH;

Access Control List Entry (ACL)



Each node has an Access Control List made up of zero or more Access Control List Entries (ACL).

The ACL block is appended to a BLK to make an access control list entry.

LITERAL

```
ACLSS_BLOCK_LENGTH = 22+BLK$K_BASE;
```

MACRO

```
SACL = BLOCK[ACLSS_BLOCK_LENGTH, BYTE] FIELD (BLKSZ_FIELDS,
                                              ACLSZ_FIELDS)
```

```
%:
```

```
FIELD ACLSZ_FIELDS =
SET
```

ACLSV_NEXT	= [0+BLK\$K_BASE, 0, 24, 0],
ACLSL_GRANT	= [3+BLK\$K_BASE, 0, 32, 0],
ACLSL_DENY	= [7+BLK\$K_BASE, 0, 32, 0],
ACLSL_BANISH	= [11+BLK\$K_BASE, 0, 32, 0],
ACLSW_UIC_MEMBER	= [15+BLK\$K_BASE, 0, 16, 0],
ACLSW_UIC_GROUP	= [17+BLK\$K_BASE, 0, 16, 0],
ACLSV_FIRST_ACLC	= [19+BLK\$K_BASE, 0, 24, 0]

```
TES;
```

LITERAL
ACLSK_BASE

= ACLSS_BLOCK_LENGTH;

Access Control List Criterion (ACLC)



Each ACL may have one or more Access Control List Criterion (ACLC) chained from it.

Each ACLC specifies one user identification criterion for the ACL entry. The user must match all the criteria for the ACL entry to apply to him.

An ACLC is appended to a BLK to form the criterion block.

LITERAL

ACLCSS_BLOCK_LENGTH = 4+BLK\$K_BASE;

MACRO

SACLC = BLOCK[ACLCSS_BLOCK_LENGTH, BYTE] FIELD (BLK\$Z_FIELDS,
ACLC\$Z_FIELDS)

%:

FIELD ACLC\$Z_FIELDS =

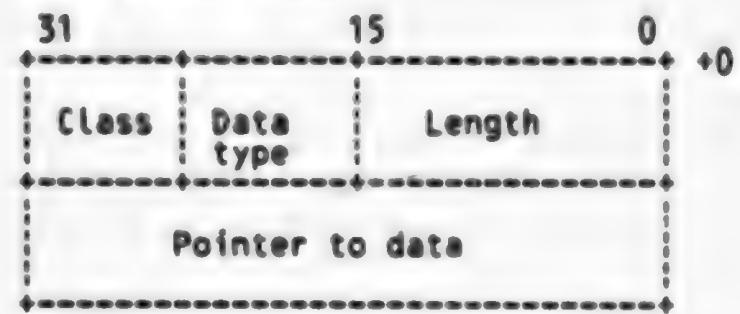
SET
ACLCV_NEXT = [0+BLK\$K_BASE, 0, 24, 0],
ACLCV_TYPE = [3+BLK\$K_BASE, 0, 8, 0],
ACLCST_STRING = [4+BLK\$K_BASE, 0, 0, 0]

TES:

LITERAL

ACLC\$K_BASE = ACLCSS_BLOCK_LENGTH;

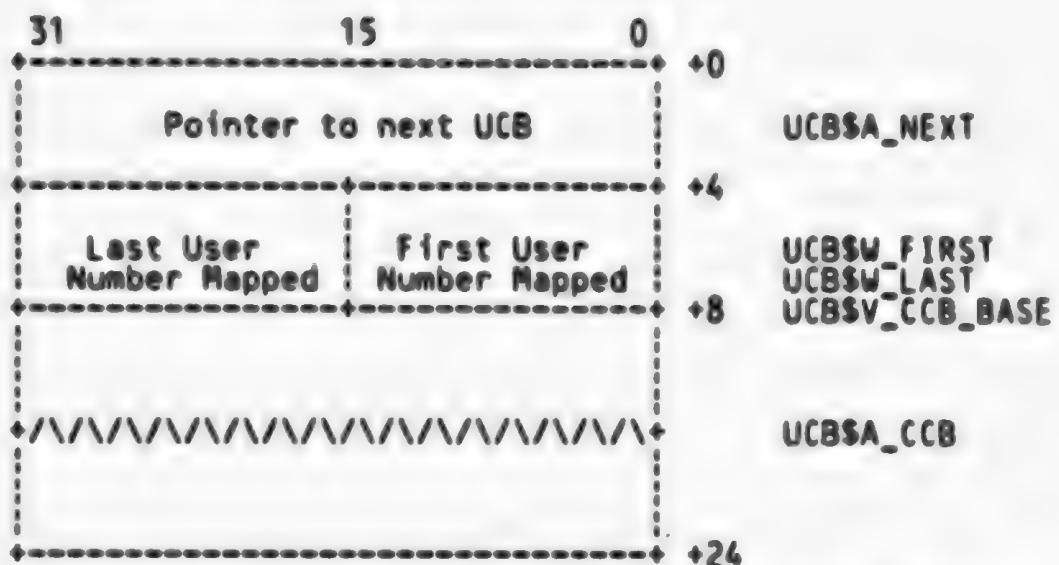
♦ String Descriptor



MACRO

```
SDSC = BLOCK[8,BYTE]  
%
```

User Control Block (UCB)



A user's context pointer is an origin 1 index into the CCBs in the UCB list. Each UCB can point to UCB\$K CCB number of CCBs. If the user passes a context number that doesn't map to an active CCB, we tell him it's an invalid context pointer.

Except in some extreme case, any image will probably never have more than 4 users active at any one time. But we can handle more!

LITERAL

UCBSK_CCB_BASE = 8,
UCBSK_CCB = 4, ! Number of CCBs/UCB
UCBSK_BLOCK_LENGTH = UCBSK_CCB_BASE + UCBSK_CCB * 4;

HACRO

SUCB = BLOCK[UCBSS_BLOCK_LENGTH, BYTE] FIELD (UCBSZ_FIELDS)
%;

FIELD UCBSZ_FIELDS =
SET

```

UCBSA_NEXT      = [0, 0, 32, 0]
UCBSW_FIRST    = [4, 0, 16, 0]
UCBSW_LAST     = [6, 0, 16, 0]
UCBSV_CCB      = [8, 0, 0, 0]
UCBSA_CCB      = [0, 0, 32, 0]

```

TES:

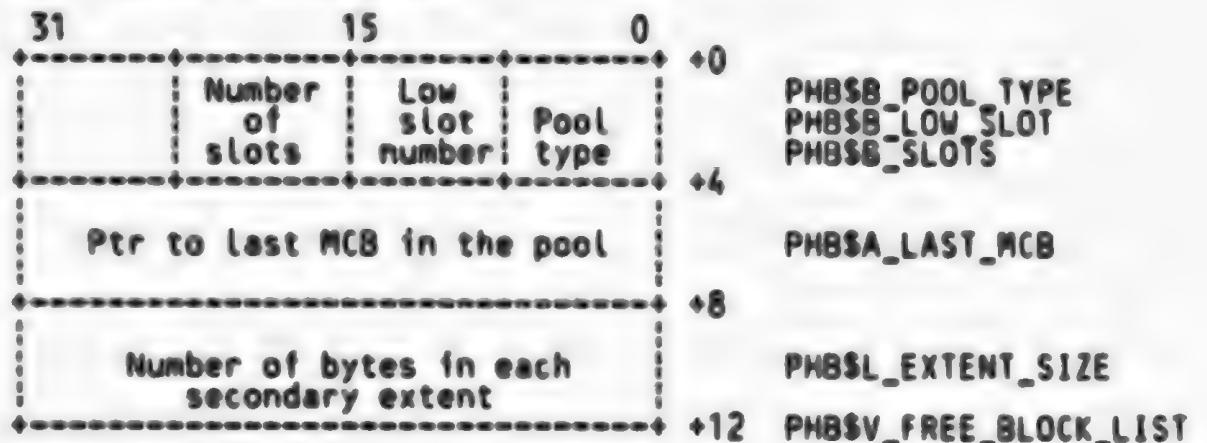
LITERAL

UCBSC_FIRST = BLOCK[0, UCBST_FIRST; UCBSS_BLOCK_LENGTH, BYTE];

STRUCTURE

UCB\$CCB(I, 0, P, S, E) = (UCB\$CCB+UCB\$K_CCB_BASE+
(I-(UCB\$CCB+UCB\$C_FIRST)<0,16,0>)*4+0)<P,S,E>;

Pool Header Block (PHB)



Pools are used to provide space for various in-core block types. Blocks that are related are usually allocated from the same pool. This is done because pools provide good locality of reference and this allocation scheme reduces page faults.

Each pool consists of one or more extents. Each extent consists of a Memory Control Block and the rest of the space allocated to the extent. Each pool has a Pool Header Block associated with it. This block identifies the MCB list, as well as contains the pool's free block list.

The free block list is a vector of linked lists. Each slot in the vector corresponds to a block type. When a block is freed, it is linked into its associated free list. When a block is requested, its free list is checked to see if it is non-empty. If so, then a block is allocated from it. Otherwise, a block is allocated from one of the pool's extents.

LITERAL

PHBSS_BLOCK_LENGTH = 12;

MACRO

SPHB = BLOCK[PHBSS_BLOCK_LENGTH, BYTE] FIELD (PHBSZ_FIELDS)
%;

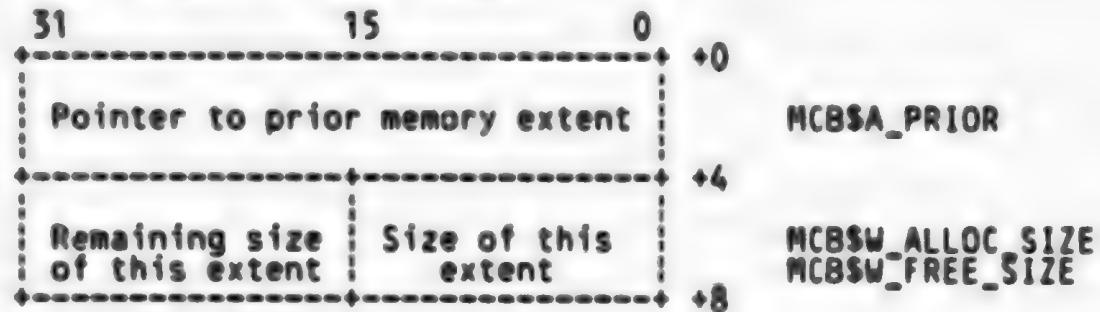
FIELD PHBSZ_FIELDS =

SET
 PHBSB_POOL_TYPE = [0, 0, 8, 0].
 PHBSB_LOW_SLOT = [1, 0, 8, 0].
 PHBSB_SLOTS = [2, 0, 8, 0].
 PHBSA_LAST_MCB = [4, 0, 32, 0].
 PHB\$L_EXTENT_SIZE = [8, 0, 32, 0].
 PHBSV_FREE_BLOCK_LIST = [12, 0, 0, 0].

TES;

1+

Memory Control Block (MCB)



Each pool has one or more Memory Control Blocks associated with it.

This list serves two purposes:

- 1) It points to each memory extent we asked LIB\$GET_VH for.
- 2) It keeps track of memory that has been allocated for the pool, but never used.

1-

LITERAL

MCBS_BLOCK_LENGTH = 8;

MACRO

SMCB = BLOCK[MCBS_BLOCK_LENGTH,BYTE] FIELD (MCBSZ_FIELDS)

%

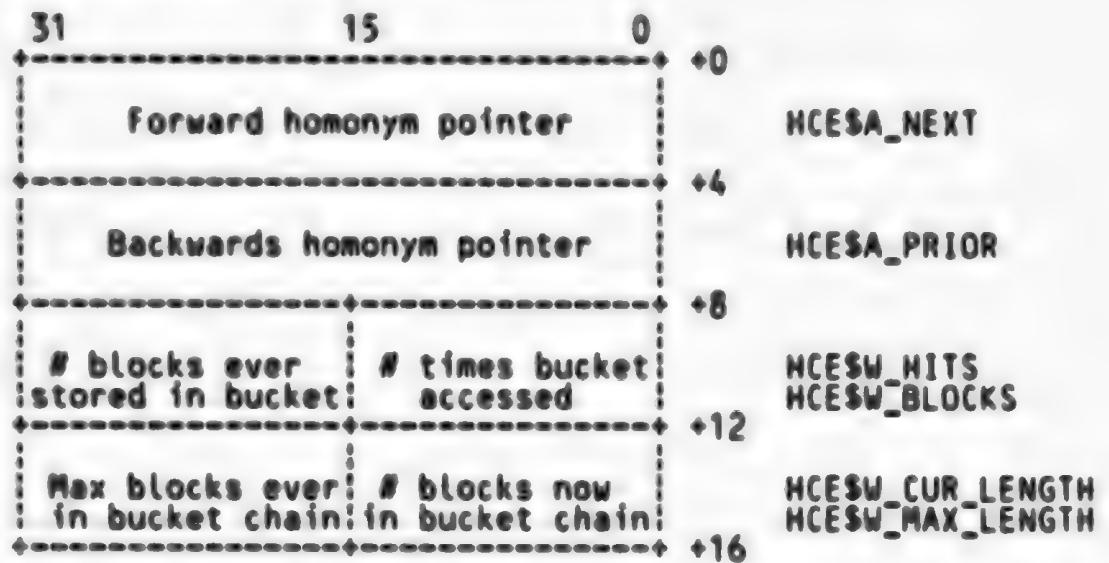
FIELD MCBSZ_FIELDS =

SET

MCBSA_PRIOR = [0, 0, 32, 0],
MCBSW_ALLOC_SIZE = [4, 0, 16, 0],
MCBSW_FREE_SIZE = [6, 0, 16, 0]

TES;

* Hash Control Entry (HCE)



HCEs are found in each user's CCB. Each HCE forms one hash bucket. We keep some statistics on the buckets in an attempt to tune the hash function.

LITERAL

```
HCESS_BLOCK_LENGTH = 16;
```

MACRO

```
$HCE = BLOCK[HCESS_BLOCK_LENGTH,BYTE] FIELD (HCESZ_FIELDS)
```

FIELD HCESZ_FIELDS =

```
SET
  HCESA_NEXT      = [0, 0, 32, 0];
  HCESA_PRIOR     = [4, 0, 32, 0];
  HCESW_HITS      = [8, 0, 16, 0];
  HCESW_BLOCKS    = [10, 0, 16, 0];
  HCESW_CUR_LENGTH = [12, 0, 16, 0];
  HCESW_MAX_LENGTH = [14, 0, 16, 0];

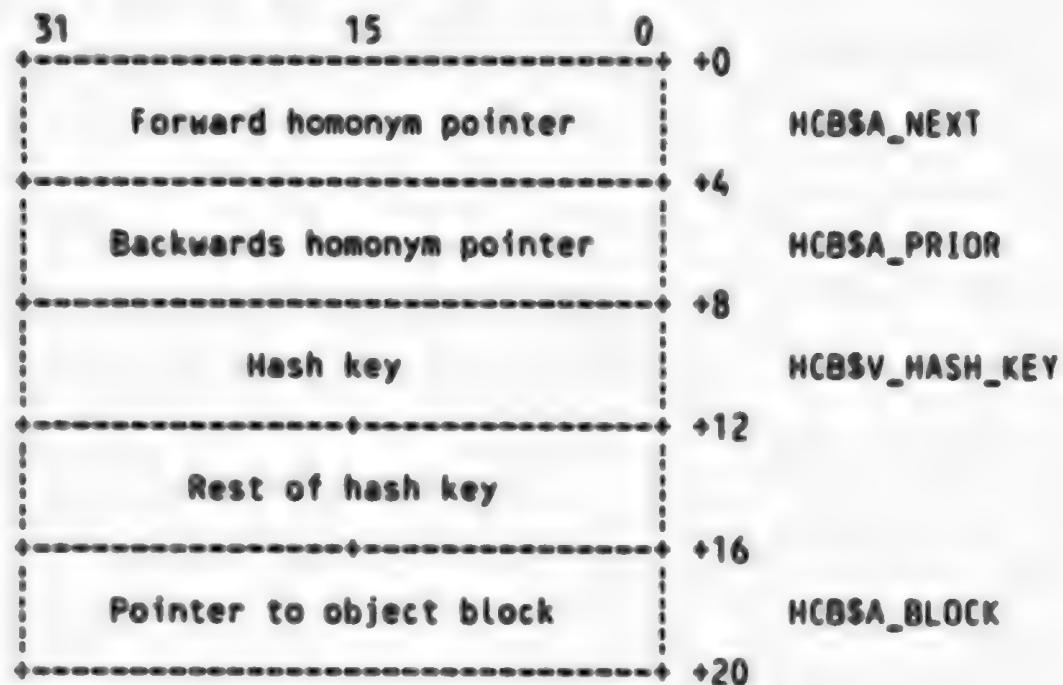
```

TES:

LITERAL

```
HCESK_HOMONYM_LIST = BLOCK[0, HCESA_NEXT:,BYTE];
```

* Hash Control Block (HCB)



LITERAL

```

HCBSK_PCB_NUMBER = 1;      ! Hash type is page number
HCBSK_LCCB_CODE  = 2;      ! Hash type is location code
HCBSK_LCCB_ADDRESS = 3;    ! Hash type is entity disk address
HCBSK_KEY_LENGTH = 8;      ! Size of a hash key
HCBS_BLOCK_LENGTH = 20;

```

MACRO

```

$HCB = BLOCK[HCBS_BLOCK_LENGTH, BYTE] FIELD (HCBSZ_FIELDS)
%;
```

FIELD HCBSZ_FIELDS =
SET

```

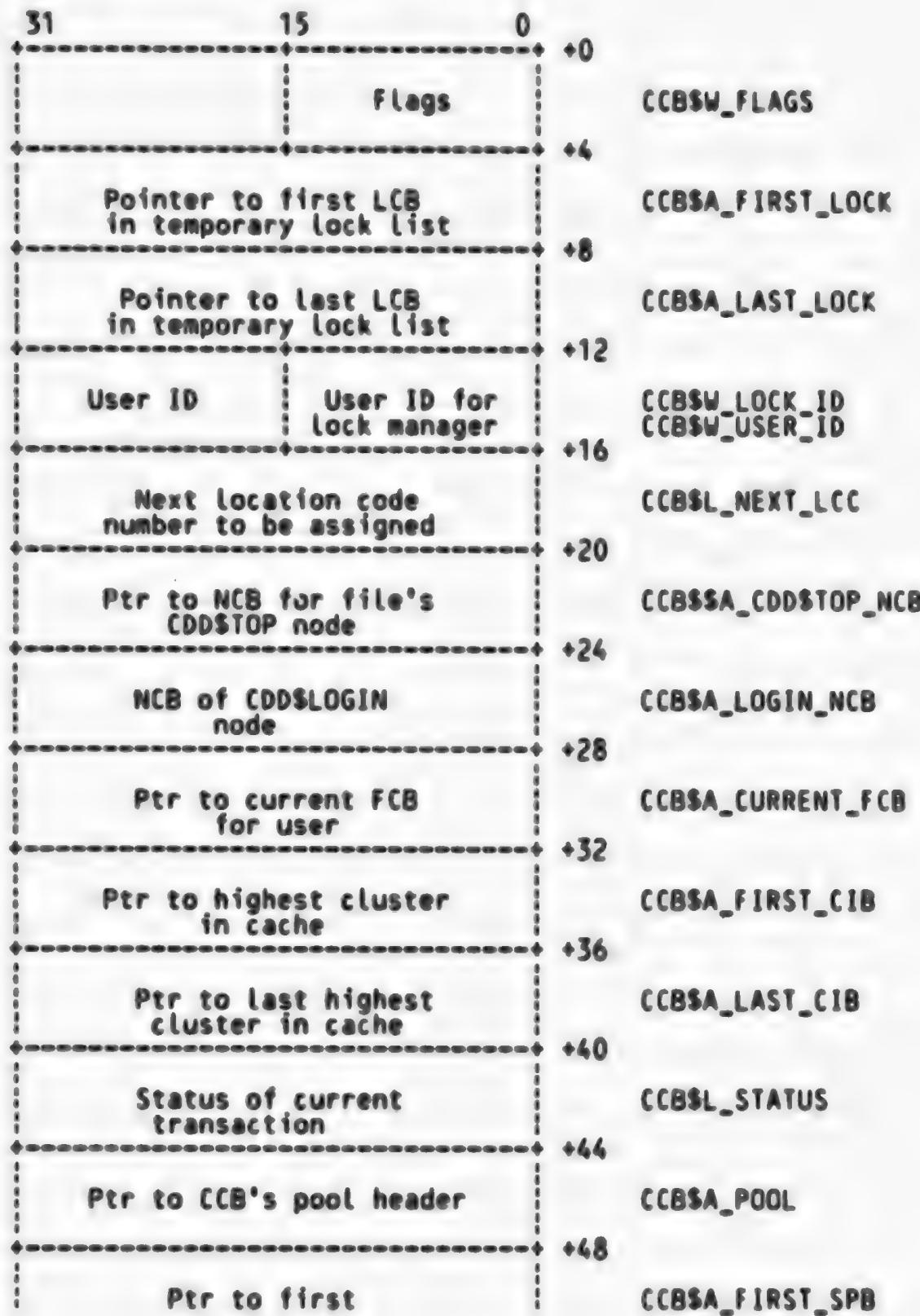
HCBSA_NEXT      = [0, 0, 32, 0];
HCBSA_PRIOR    = [4, 0, 32, 0];
HCBSV_HASH_KEY = [8, 0, 0, 0];
HCBSA_BLOCK    = [16, 0, 32, 0]

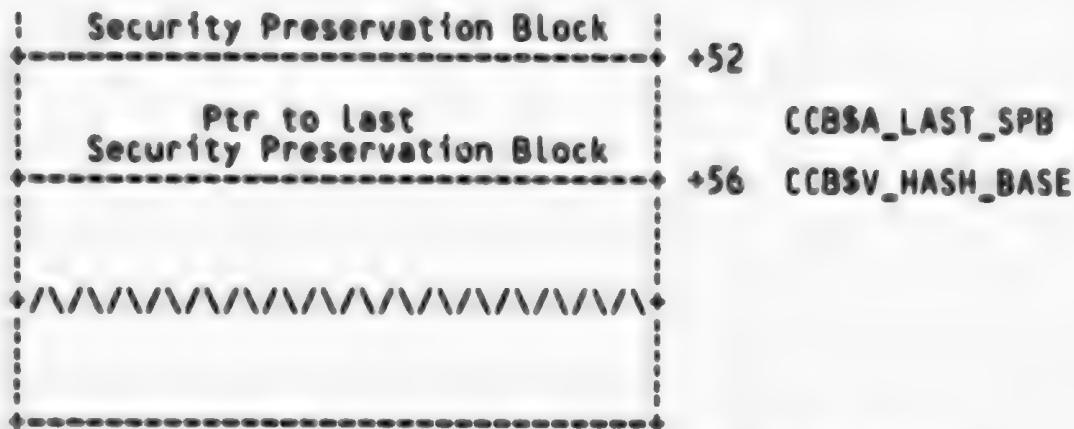
```

TES:

LITERAL
HCBSK_HOMONYM_LIST = BLOCK[0, HCBSA_NEXT; , BYTE];

Context Control Block (CCB)





The hash table consists of a number of hash entries.
See the HCE description for the format of these entries.

LITERAL

```
CCBSK_HASH_TABLE_SIZE = 151.
CCBSS_BLOCK_LENGTH    = 56 + (CCBSK_HASH_TABLE_SIZE * HCESS_BLOCK_LENGTH);
```

MACRO

```
$CCB = BLOCK[CCBSS_BLOCK_LENGTH,BYTE] FIELD (CCBSZ_FIELDS)
%;
```

FIELD CCBSZ_FIELDS =

```
SET
  CCBSU_FLAGS          = [0, 0, 16, 0].
  CCBSV_CORRUPT        = [0, 0, 16, 0].
  CCBSA_FIRST_LOCK     = [4, 0, 32, 0].
  CCBSA_LAST_LOCK      = [8, 0, 32, 0].
  CCBSW_LOCK_ID        = [12, 0, 16, 0].
  CCBSW_USER_ID        = [14, 0, 16, 0].
  CCBSL_NEXT_LCC       = [16, 0, 32, 0].
  CCBSA_CDDSTOP_NCB   = [20, 0, 32, 0].
  CCBSA_LOGIN_NCB     = [24, 0, 32, 0].
  CCBSA_CURRENT_FCB   = [28, 0, 32, 0].
  CCBSA_FIRST_CIB      = [32, 0, 32, 0].
  CCBSA_LAST_CIB       = [36, 0, 32, 0].
  CCBSL_STATUS         = [40, 0, 32, 0].
  CCBSA_POOL           = [44, 0, 32, 0].
  CCBSA_FIRST_SPB      = [48, 0, 32, 0].
  CCBSA_LAST_SPB       = [52, 0, 32, 0].
  CCBSV_HASH_BASE      = [56, 0, 0, 0]

  ! Stream is corrupt
```

TES:

LITERAL

```
CCBSK_LOCK_LIST      = BLOCK[0, CCBSA_FIRST_LOCK:, BYTE].
CCBSK_CLUSTER_LIST   = BLOCK[0, CCBSA_FIRST_CIB:, BYTE].
CCBSS_HASH_BASE      = BLOCK[0, CCBSV_HASH_BASE:, BYTE].
CCBSK_SPB_LIST       = BLOCK[0, CCBSA_FIRST_SPB:, BYTE];
```

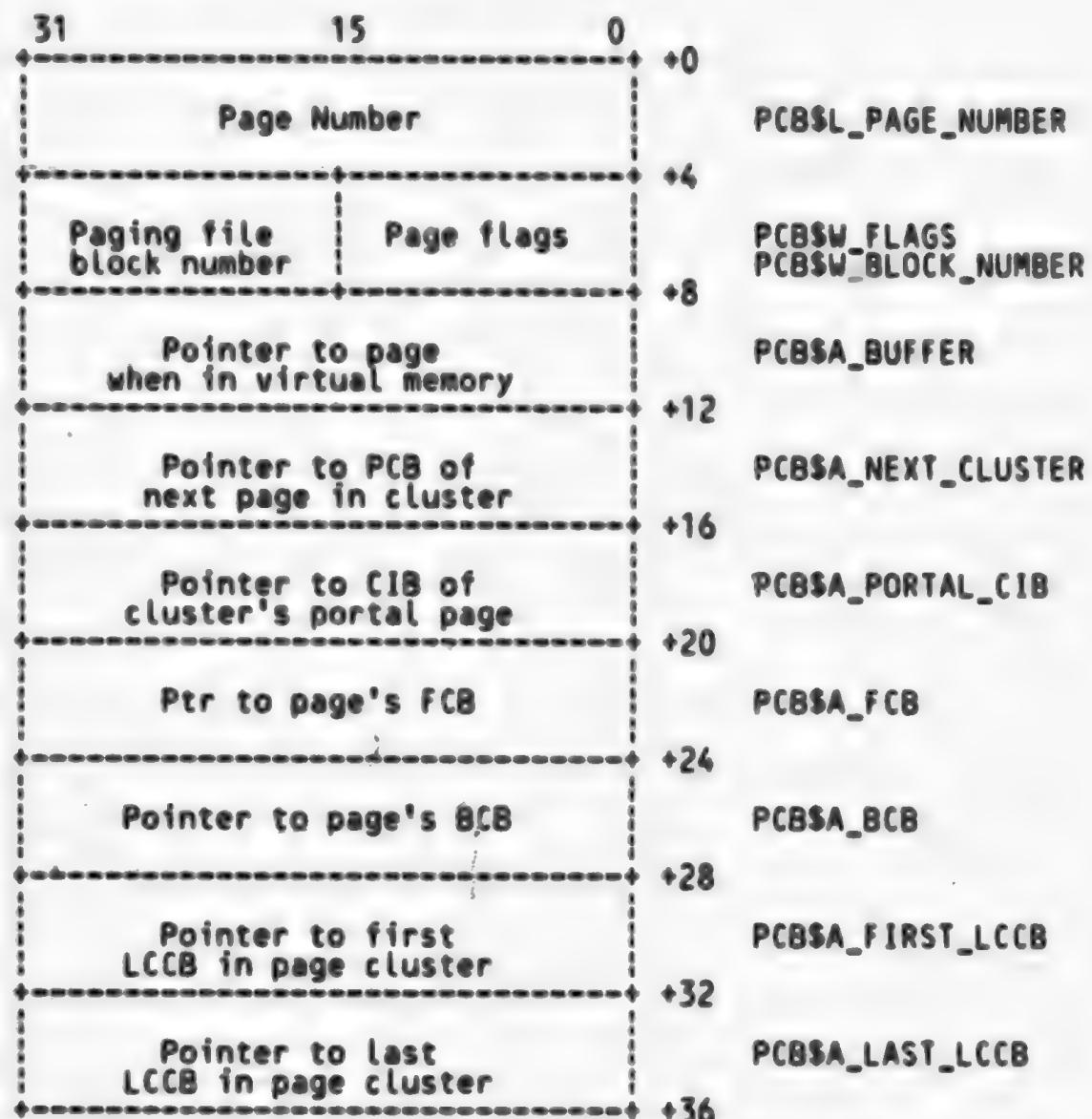
STRUCTURE

```
[CBSHASH[1, 0, P, S, E] =
```

(CCB\$HASH+CCBSS_HASH_BASE+0+(I*HCESS_BLOCK_LENGTH))<P,S,E>:

LITERAL
CCBSM_CORRUPT = 1^1 - 1^0: ! Stream is corrupt

Page Control Block (PCB)



A Page Control Block (PCB) exists for every page in the cache, and for pages that only have presence locks on them (thus they're not in the staging cache). Portal pages have a Cluster Information Block (CIB) appending to their PCB.

LITERAL
PCBSS_BLOCK_LENGTH = 36;

MACRO
SPCB = BLOCK[PCBSS_BLOCK_LENGTH,BYTE] FIELD (PCBSZ_FIELDS)
%;

FIELD SET PCB\$Z_FIELDS =

PCBSL_PAGE_NUMBER	= [0, 0, 32, 0],	
PCBSW_FLAGS	= [4, 0, 16, 0],	
PCBSV_MODIFIED	= [4, 0, 1, 0],	Must write back to dict.
PCBSV_NEW_PAGE	= [4, 3, 1, 0],	Page was in free chain
PCBSV_READ_ONLY	= [4, 4, 1, 0],	Page is read only
PCBSV_FREE_PAGE	= [4, 5, 1, 0],	Page is a free page
PCBSV_AVAILABLE	= [4, 6, 1, 0],	Page available in cache
PCBSV_PORTAL_PAGE	= [4, 7, 1, 0],	Page is portal page, [IB follows
PCBSW_BLOCK_NUMBER	= [6, 0, 16, 0],	
PCBSA_BUFFER	= [8, 0, 32, 0],	
PCBSA_NEXT_CLUSTER	= [12, 0, 32, 0],	
PCBSA_PORTAL_CIB	= [16, 0, 32, 0],	
PCBSA_FCB	= [20, 0, 32, 0],	
PCBSA_BCB	= [24, 0, 32, 0],	
PCBSA_FIRST_LCCB	= [28, 0, 32, 0],	
PCBSA_LAST_LCCB	= [32, 0, 32, 0]	

TES:

LITERAL
PCBSK_LCCB_LIST

= BLOCK[0, PCBSA_FIRST_LCCB: , BYTE];

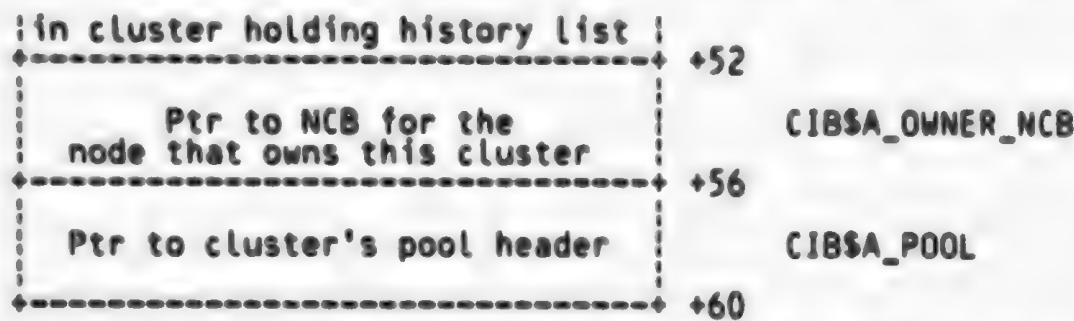
LITERAL

PCBSM_MODIFIED	= 1^1 - 1^0,	Must write back to dict.
PCBSM_NEW_PAGE	= 1^4 - 1^3,	Page was in free chain
PCBSM_READ_ONLY	= 1^5 - 1^4,	Page is read only
PCBSM_FREE_PAGE	= 1^6 - 1^5,	Page is a free page
PCBSM_AVAILABLE	= 1^7 - 1^6,	Page available in cache
PCBSM_PORTAL_PAGE	= 1^8 - 1^7:	Portal page, [IB follows

Cluster Information Block (CIB)

31	15	0	+0
		Cluster flags	+0
Retrieval lock ref count	Presence lock ref count		+4
Delete lock ref count	Update lock ref count		+8
Pointer to first deleted child cluster			+12
Pointer to last deleted child cluster			+16
Ptr to newest lock granted for this page			+20
Ptr to oldest lock granted for this page			+24
Ptr to next cluster at this directory level			+28
Ptr to prior cluster at this directory level			+32
Ptr to first cluster at next directory level			+36
Ptr to last cluster at next directory level			+40
Ptr to CIB in next highest directory level			+44
Ptr to portal CIB			+48

CIBSW_FLAGS
 CIBSV_REF_COUNTS
 CIBSU_PRESENCE_COUNT
 CIBSW_RETRIEVAL_COUNT
 CIBSW_UPDATE_COUNT
 CIBSW_DELETE_COUNT
 CIBSA_FIRST_DELETED_CIB
 CIBSA_LAST_DELETED_CIB
 CIBSA_LAST_LCB
 CIBSA_FIRST_LCB
 CIBSA_NEXT_SIBLING
 CIBSA_PRIOR_SIBLING
 CIBSA_FIRST_CHILD
 CIBSA_LAST_CHILD
 CIBSA_PARENT_CIB
 CIBSA_HISTORY_CIB



Each clusters' portal page has a CIB associated with it.
 This block is physically appended to the PCB of the cluster's
 portal page.

The CIB also points to the cluster's pool header.

The CIB\$V_REF_COUNTS lists must have each of the lock ref counts
 in the same order as the LOCK\$K_xxx lock request symbols.

!-

LITERAL

```
CIBSS_BLOCK_LENGTH = 60 + PCBSS_BLOCK_LENGTH;
```

MACRO

```
SCIB = BLOCK[CIBSS_BLOCK_LENGTH,BYTE] FIELD {CIBSZ_FIELDS,
                                              PCBSZ_FIELDS}
```

%:

FIELD CIBSZ_FIELDS =

```

SET
CIBSW_FLAGS          = [0+PCBSS_BLOCK_LENGTH, 0, 16, 0],
CIBSV_LOCKED_DELETE  = [0+PCBSS_BLOCK_LENGTH, 0, 1, 0],
CIBSV_LOCKED_UPDATE  = [0+PCBSS_BLOCK_LENGTH, 1, 1, 0],
CIBSV_LOCKED_RETRIEVAL = [0+PCBSS_BLOCK_LENGTH, 2, 1, 0],
CIBSV_LOCKED_PRESENCE = [0+PCBSS_BLOCK_LENGTH, 3, 1, 0],
CIBSV_LOCKED          = [0+PCBSS_BLOCK_LENGTH, 0, 4, 0],
CIBSV_LOGIN           = [0+PCBSS_BLOCK_LENGTH, 4, 1, 0],
CIBSV_COMPLETE         = [0+PCBSS_BLOCK_LENGTH, 5, 1, 0],
CIBSV_NEW              = [0+PCBSS_BLOCK_LENGTH, 6, 1, 0],
CIBSV_HISTORY          = [0+PCBSS_BLOCK_LENGTH, 7, 1, 0],
CIBSV_REF_COUNTS       = [4+PCBSS_BLOCK_LENGTH, 0, 0, 0],
CIBSW_PRESENCE_COUNT  = [4+PCBSS_BLOCK_LENGTH, 0, 16, 0],
CIBSW_RETRIEVAL_COUNT = [6+PCBSS_BLOCK_LENGTH, 0, 16, 0],
CIBSW_UPDATE_COUNT    = [8+PCBSS_BLOCK_LENGTH, 0, 16, 0],
CIBSW_DELETE_COUNT    = [10+PCBSS_BLOCK_LENGTH, 0, 16, 0],
CIBSA_FIRST_DELETED_CIB = [12+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_LAST_DELETED_CIB = [16+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_LAST_LCB          = [20+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_FIRST_LCB         = [24+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_NEXT_SIBLING      = [28+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_PRIOR_SIBLING     = [32+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_FIRST_CHILD        = [36+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_LAST_CHILD         = [40+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_PARENT_CIB         = [44+PCBSS_BLOCK_LENGTH, 0, 32, 0],

```

CIBSA_HISTORY [10] = [48+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_OWNER_NCB = [52+PCBSS_BLOCK_LENGTH, 0, 32, 0],
CIBSA_POOL = [56+PCBSS_BLOCK_LENGTH, 0, 32, 0]
TES;

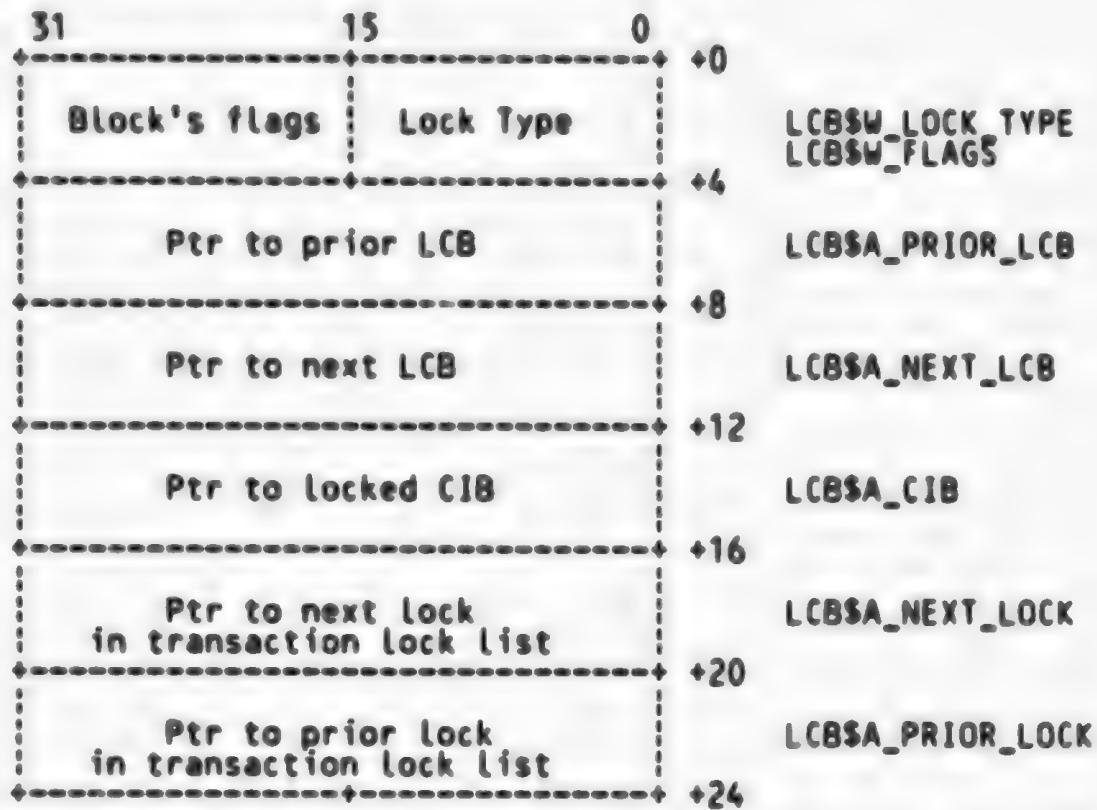
LITERAL

CIBSK_LOCK_LIST = BLOCK[0, CIBSA_LAST_LCB:, BYTE],
CIBSK_SIBLING_LIST = BLOCK[0, CIBSA_NEXT_SIBLING:, BYTE],
CIBSK_CHILD_LIST = BLOCK[0, CIBSA_FIRST_CHILD:, BYTE];

LITERAL

CIBSM_LOCKED_DELETE = 1^1 - 1^0,
CIBSM_LOCKED_UPDATE = 1^2 - 1^1,
CIBSM_LOCKED_RETRIEVAL = 1^3 - 1^2,
CIBSM_LOCKED_PRESENCE = 1^4 - 1^3,
CIBSM_LOCKED = 1^4 - 1^0,
CIBSM_LOGIN = 1^5 - 1^4, ! Cluster is on login path
CIBSM_COMPLETE = 1^6 - 1^5, ! Cluster is complete in cache
CIBSM_NEW = 1^7 - 1^6, ! Cluster read in this transaction
CIBSM_HISTORY = 1^8 - 1^7; ! History list is in cluster

Lock Control Block (LCB)



LCBs are used to keep track of which locks exist on a cluster (CIB).

Each LCB is linked to its CIB, and to other LCBs for that cluster. When an LCB is granted, it is placed in the transaction's lock list until the transaction terminates.

LITERAL

```
LCBSS_BLOCK_LENGTH = 24;
```

MACRO

```
$LCB = BLOCK[LCBSS_BLOCK_LENGTH,BYTE] FIELD (LCBSZ_FIELDS)
```

FIELD

```
LCBSZ_FIELDS =
```

LCBSW_LOCK_TYPE	= [0, 0, 16, 0],
LCBSW_FLAGS	= [2, 0, 16, 0],
LCBSV_CURRENT	= [2, 1, 1, 0] : Allocated in current transaction
LCBSA_PRIOR_LCB	= [4, 0, 32, 0],
LCBSA_NEXT_LCB	= [8, 0, 32, 0],
LCBSA_CIB	= [12, 0, 32, 0],
LCBSA_NEXT_LOCK	= [16, 0, 32, 0],
LCBSA_PRIOR_LOCK	= [20, 0, 32, 0]

```
TES;
```

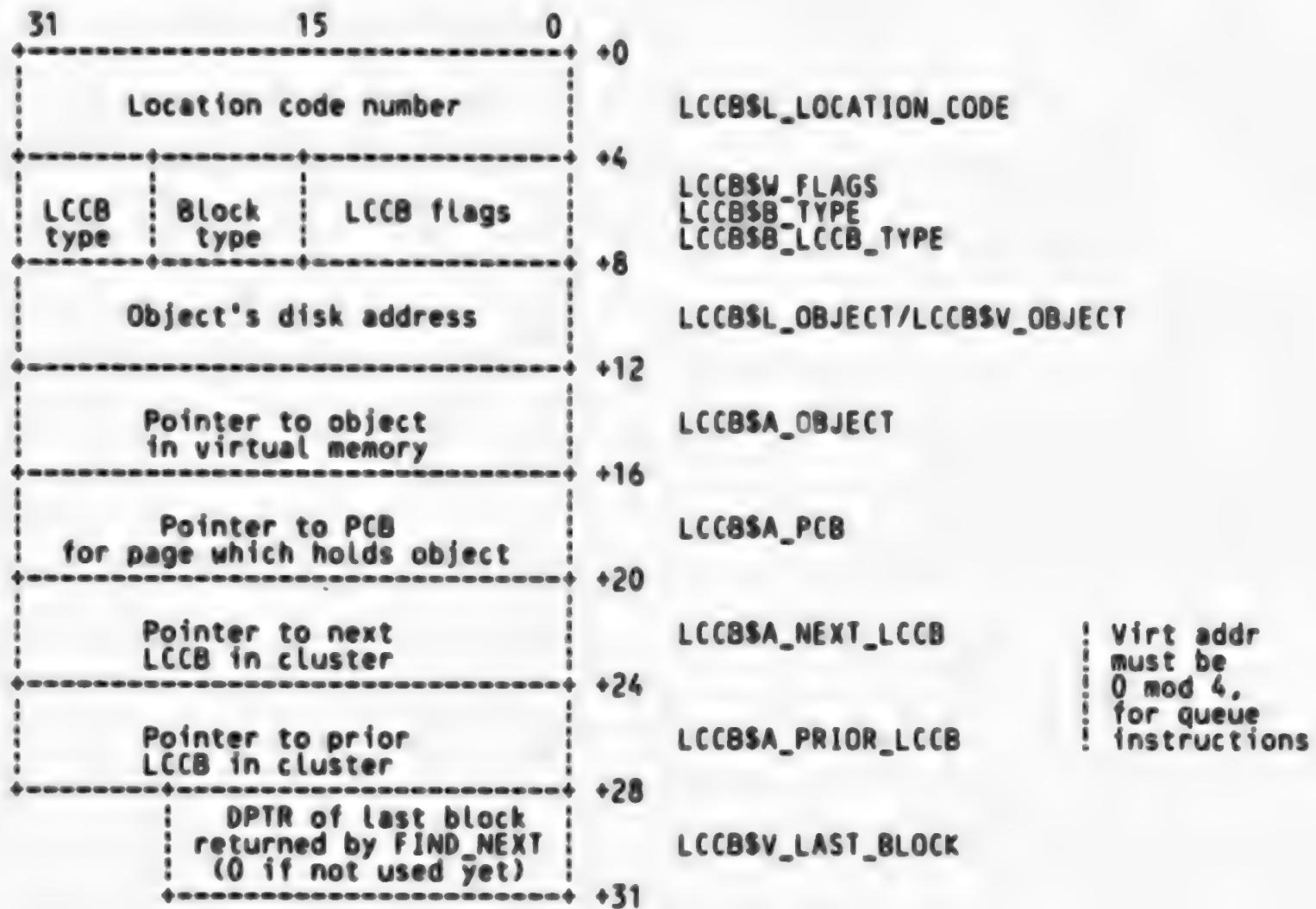
LITERAL
LCBSK_LOCK_LIST
LCBSK_TEMP_LOCK_LIST

= BLOCK[0, LCB8A_PRIOR_LCB; , BYTE];
= BLOCK[0, LCB8A_NEXT_LOCK; , BYTE];

LITERAL
LCBSM_CURRENT

= 1^2 - 1^1; ! Allocated in current transaction

* Location Code Control Block (LCCB)



! Virt addr
 must be
 0 mod 4,
 for queue
 instructions

Location codes are used to provide a convenient means for the user to identify a particular object. Each location code is associated with an LCCB. The following objects may be assigned location code:

- 1) Entities
- 2) Lists
- 3) Nodes

LITERAL
LCCBSS_BLOCK_LENGTH = 31;

MACRO

SLCCB = BLOCK[LCCBSZ_BLOCK_LENGTH, BYTE] FIELD (LCCBSZ_FIELDS)
%;

FIELD LCCBSZ_FIELDS =
SET

LCCBSL_LOCATION_CODE	= [0, 0, 32, 0],
LCCBSW_FLAGS	= [4, 0, 16, 0],
LCCBSV_AVAILABLE	= [4, 0, 1, 0],
LCCBSV_MUST_SCAN	= [4, 1, 1, 0],
LCCBSV_GHOST	= [4, 2, 1, 0],
LCCBSV_DIRECTORY	= [4, 3, 1, 0],
LCCBSV_TERMINAL	= [4, 4, 1, 0],
LCCBSV_ENTITY_ATT	= [4, 5, 1, 0],
LCCBSV_ENTITY_LIST_ATT	= [4, 6, 1, 0],
LCCBSV_ENTITY_LIST	= [4, 7, 1, 0],
LCCBSV_STRING_LIST_ATT	= [4, 8, 1, 0],
LCCBSB_TYPE	= [6, 0, 8, 0],
LCCBSB_LCCB_TYPE	= [7, 0, 8, 0],
LCCBSL_OBJECT	= [8, 0, 32, 0],
LCCBSV_OBJECT	= [8, 0, 24, 0],
LCCBSA_OBJECT	= [12, 0, 32, 0],
LCCBSA_PCB	= [16, 0, 32, 0],
LCCBSA_NEXT_LCCB	= [20, 0, 32, 0],
LCCBSA_PRIOR_LCCB	= [24, 0, 32, 0],
LCCBSV_LAST_BLOCK	= [28, 0, 24, 0]

TES:

LITERAL
LCCBSK_LCCB_LIST

= BLOCK[0, LCCBSA_NEXT_LCCB; , BYTE];

LITERAL

LCCBSM_AVAILABLE	= 1^1 - 1^0,	Object's virtual addr is known
LCCBSM_MUST_SCAN	= 1^2 - 1^1,	Protection tree must be scanned
LCCBSM_GHOST	= 1^3 - 1^2,	LCCB may not be fetched by LCC
LCCBSM_DIRECTORY	= 1^4 - 1^3,	Directory NCB
LCCBSM_TERMINAL	= 1^5 - 1^4,	Terminal NCB
LCCBSM_ENTITY_ATT	= 1^6 - 1^5,	Entity attribute LCCB
LCCBSM_ENTITY_LIST_ATT	= 1^7 - 1^6,	Entity list attribute LCCB
LCCBSM_ENTITY_LIST	= 1^8 - 1^7,	Entity list ECCB
LCCBSM_STRING_LIST_ATT	= 1^9 - 1^8,	String list LCCB

LCCBSK_LCCB_TYPE_FIRST

= 1,

| LCCB includes NCB

LCCBSK_NCB

= 1,

| LCCB includes ECCB

LCCBSK_ECCB

= 2,

| LCCB

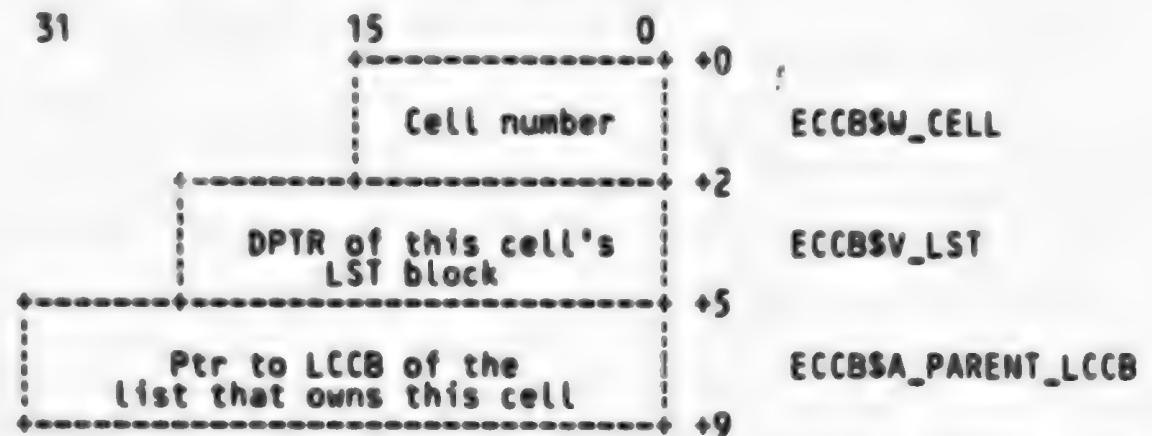
LCCBSK_LCCB

= 3,

| LCCB

LCCBSK_LCCB_TYPE_LAST

Entity Cell Control Block (ECCB)



Each cell in an entity list may be assigned a location code.

This block is appended to the location code's LCCB to name the specific cell represented by the location code.

LITERAL

ECCBSS_BLOCK_LENGTH = LCCBSS_BLOCK_LENGTH + 9;

MACRO

\$ECCB = BLOCK[ECCBSS_BLOCK_LENGTH,BYTE] FIELD (ECCBSZ_FIELDS,
LCCBSZ_FIELDS)

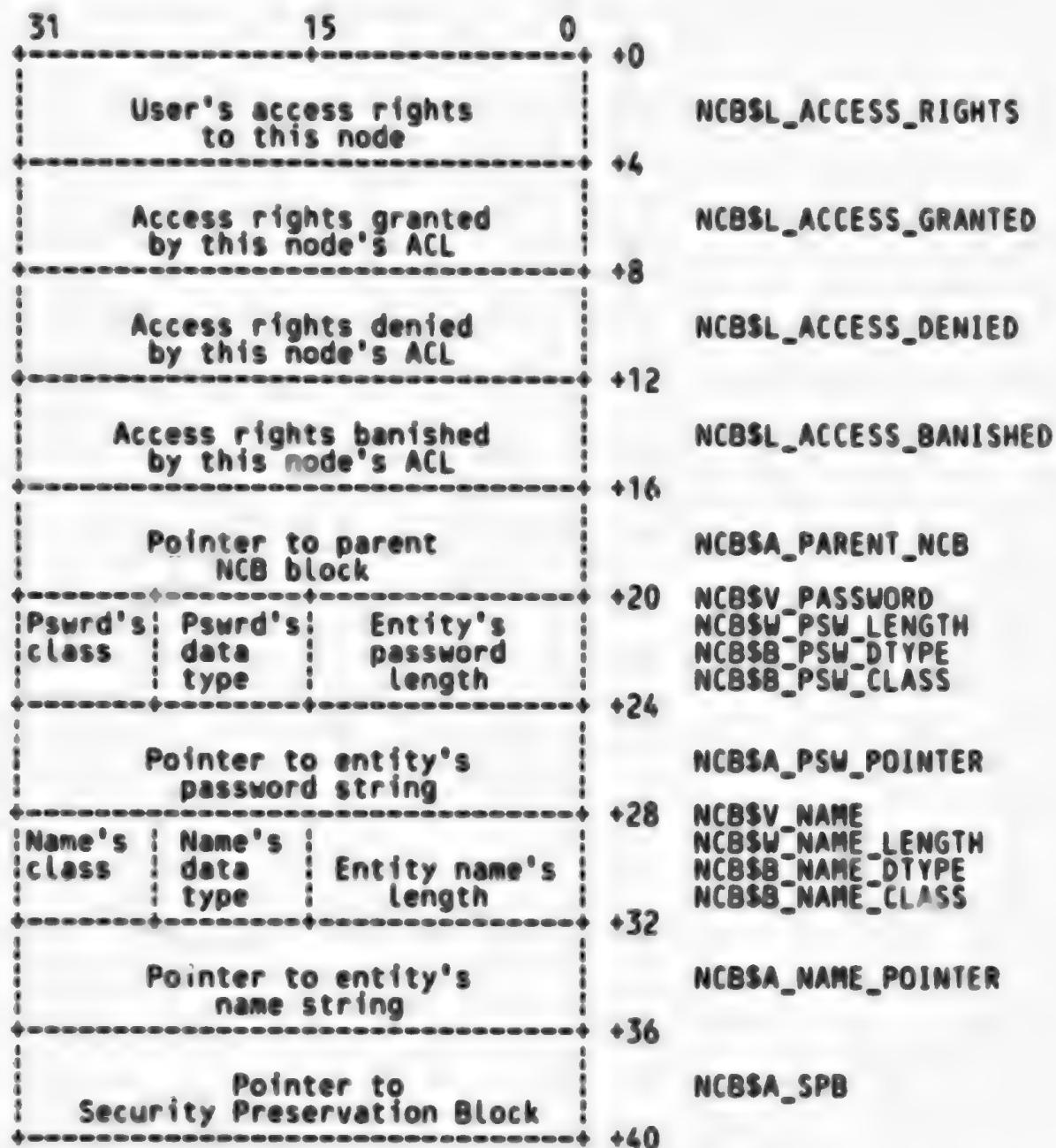
%;

FIELD ECCBSZ_FIELDS =

SET
ECCBSW_CELL = [0+LCCBSS_BLOCK_LENGTH, 0, 16, 0].
ECCBSV_LST = [2+LCCBSS_BLOCK_LENGTH, 0, 24, 0].
ECCBSA_PARENT_LCCB = [5+LCCBSS_BLOCK_LENGTH, 0, 32, 0]

TES;

Node Control Block (NCB)



The LCCB for a node is followed by an NCB. This gives additional information needed for the node.

LITERAL

NCBSS_BLOCK_LENGTH = 40+LCCBSS_BLOCK_LENGTH;

MACRO

\$NCB = BLOCK[NCBSS_BLOCK_LENGTH, BYTE] FIELD (LCCBSS_FIELDS,
NCBSS_FIELDS)

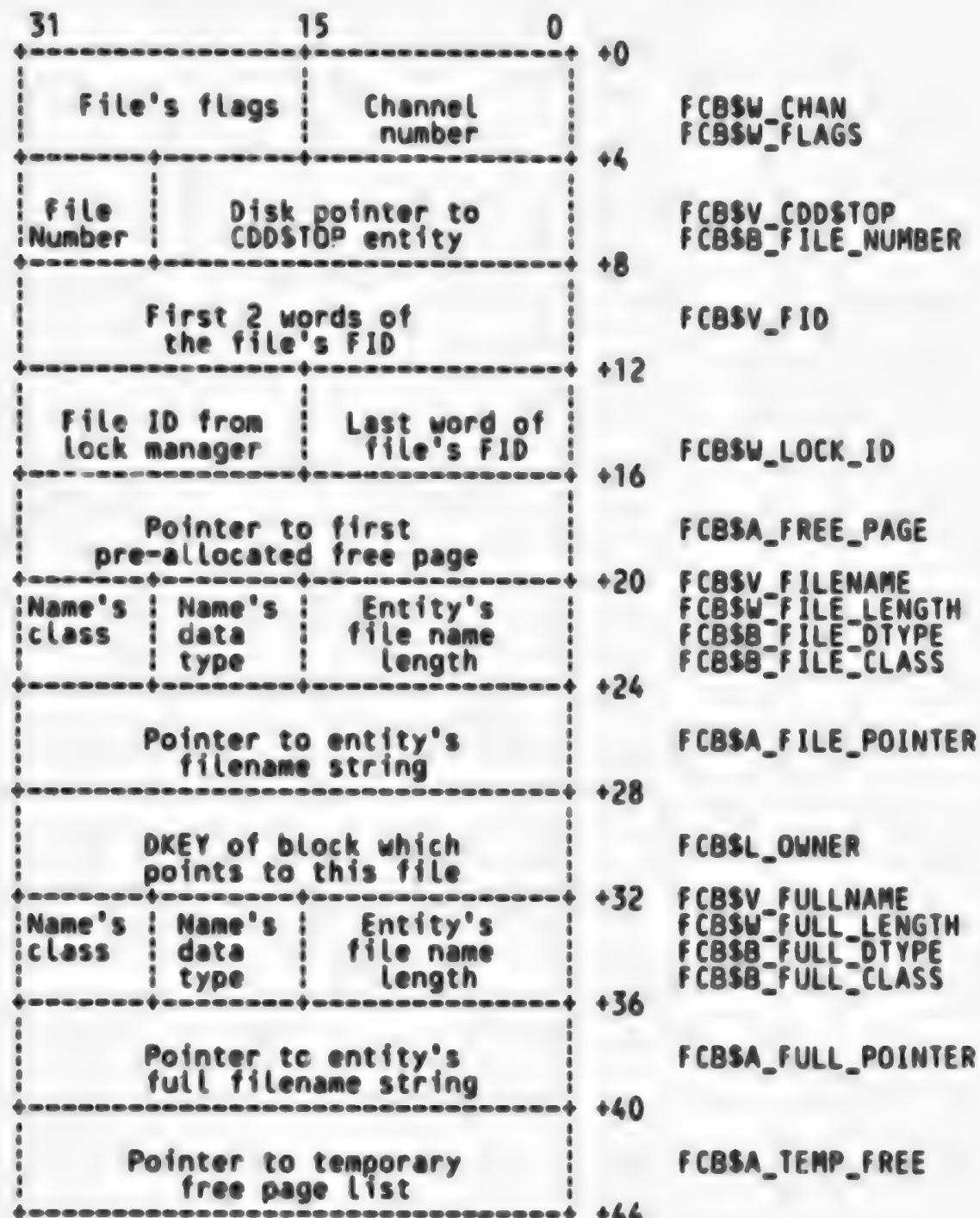
%:

FIELD NCBSS_FIELDS =
SET

NCBSL_ACCESS_RIGHTS	= [0+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSL_ACCESS_GRANTED	= [4+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSL_ACCESS_DENIED	= [8+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSL_ACCESS_BANISHED	= [12+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSA_PARENT_NCB	= [16+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSV_PASSWORD	= [20+LCCBSS_BLOCK_LENGTH, 0, 0, 0],
NCBSW_PSW_LENGTH	= [20+LCCBSS_BLOCK_LENGTH, 0, 16, 0],
NCBSB_PSW_DTYPE	= [22+LCCBSS_BLOCK_LENGTH, 0, 8, 0],
NCBSB_PSW_CLASS	= [23+LCCBSS_BLOCK_LENGTH, 0, 8, 0],
NCBSA_PSW_POINTER	= [24+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSV_NAME	= [28+LCCBSS_BLOCK_LENGTH, 0, 0, 0],
NCBSW_NAME_LENGTH	= [28+LCCBSS_BLOCK_LENGTH, 0, 16, 0],
NCBSB_NAME_DTYPE	= [30+LCCBSS_BLOCK_LENGTH, 0, 8, 0],
NCBSB_NAME_CLASS	= [31+LCCBSS_BLOCK_LENGTH, 0, 8, 0],
NCBSA_NAME_POINTER	= [32+LCCBSS_BLOCK_LENGTH, 0, 32, 0],
NCBSA_SPB	= [36+LCCBSS_BLOCK_LENGTH, 0, 32, 0]

TES:

File Control Block (FCB)

FCBSW_CHAN
FCBSW_FLAGSFCBSV_CDDSTOP
FCBSB_FILE_NUMBER

FCBSV_FID

FCBSW_LOCK_ID

FCBSA_FREE_PAGE

FCBSV_FILENAME
FCBSW_FILE_LENGTH
FCBSB_FILE_DTYPE
FCBSB_FILE_CLASS

FCBSA_FILE_POINTER

FCBSL_OWNER

FCBSV_FULLNAME
FCBSW_FULL_LENGTH
FCBSB_FULL_DTYPE
FCBSB_FULL_CLASS

FCBSA_FULL_POINTER

FCBSA_TEMP_FREE

Each active (open) dictionary file has an FCB.

The FCBSL_OWNER field holds the DKEY of the attribute block (FIL) which pointed to the file. The primary dictionary file has key 0, while files that are not currently pointed to have a value of -1

in the owner field.

LITERAL

FCBSK_FILE_LIMIT = 255.
FCBSS_BLOCK_LENGTH = 44;

MACRO

\$FCB = BLOCK[FCBSS_BLOCK_LENGTH,BYTE] FIELD (FCBSZ_FIELDS)

FIELD SET FCBSZ_FIELDS =

FCBSW_CHAN	= [0, 0, 16, 0]
FCBSW_FLAGS	= [2, 0, 16, 0]
FCBSV_READ_ONLY	= [2, 0, 1, 0]
FCBSV_ROOT	= [2, 1, 1, 0]
FCBSV_CDDSTOP	= [4, 0, 24, 0]
FCBSB_FILE_NUMBER	= [7, 0, 8, 0]
FCBSV_FID	= [8, 0, 0, 0]
FCBSW_LOCK_ID	= [14, 0, 16, 0]
FCBSA_FREE_PAGE	= [16, 0, 32, 0]
FCBSV_FILENAME	= [20, 0, 0, 0]
FCBSW_FILE_LENGTH	= [20, 0, 16, 0]
FCBSB_FILE_DTYPE	= [22, 0, 8, 0]
FCBSB_FILE_CLASS	= [23, 0, 8, 0]
FCBSA_FILE_POINTER	= [24, 0, 32, 0]
FCBSL_OWNER	= [28, 0, 32, 0]
FCBSV_FULLNAME	= [32, 0, 0, 0]
FCBSW_FULL_LENGTH	= [32, 0, 16, 0]
FCBSB_FULL_DTYPE	= [34, 0, 8, 0]
FCBSB_FULL_CLASS	= [35, 0, 8, 0]
FCBSA_FULL_POINTER	= [36, 0, 32, 0]
FCBSA_TEMP_FREE	= [40, 0, 32, 0]

TES:

LITERAL

FCBSM_READ_ONLY = 1^1 - 1^0. ! File can only be opened for read
FCBSM_ROOT = 1^2 - 1^1. ! FCB is the root dictionary file

♦ Pre-Allocated Page Block (PAPB)



Pre-allocated pages are represented by PAPBs linked onto the FCB in which the pages reside. Each PAPB has the page number of the free page it represents.

LITERAL

```
PAPBSS_BLOCK_LENGTH = 8;
```

MACRO

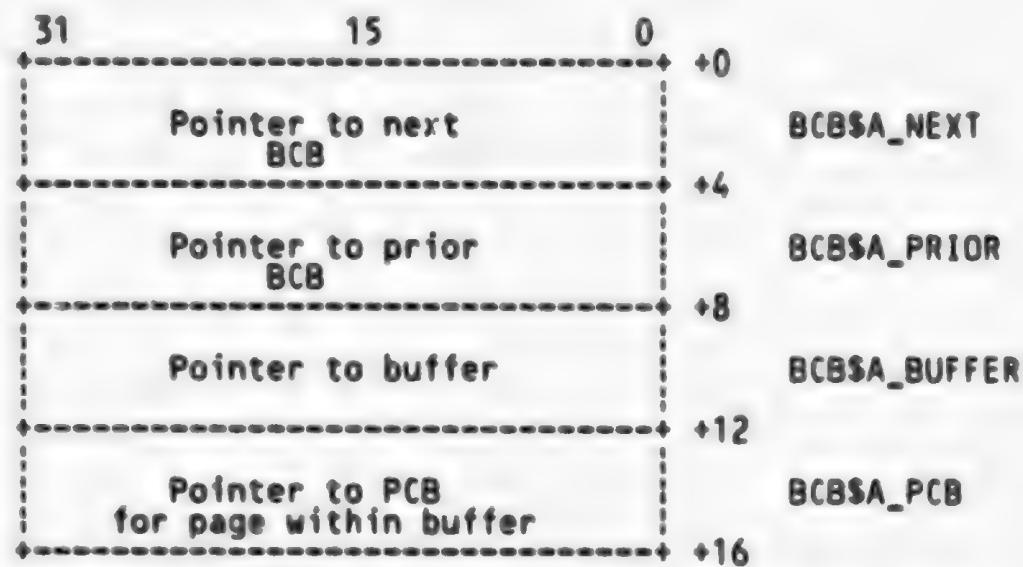
```
SPAPB = BLOCK[PAPBSS_BLOCK_LENGTH,BYTE] FIELD (PAPBSZ_FIELDS)
```

```
%;
```

FIELD PAPBSZ_FIELDS =

```
SET PAPBSA_NEXT = [0, 0, 32, 0],  
PAPBSL_NUMBER = [4, 0, 32, 0]  
TES;
```

Buffer Control Block (BCB)



The buffer control blocks (BCBs) are used to control the pages that are in memory.

The BCBs are linked into a queue. When a buffer is needed, the last buffer in the queue is assigned.

The associated PCB contains information about the page and the buffer. If a page is modified while in the buffer, the page is written back to the work file before the buffer is reused.

LITERAL

```
BCBSK_NUMBER      = 16.    ! Number of in-core buffers
BCBSS_BLOCK_LENGTH = 16;
```

MACRO

```
$BCB = BLOCK[BCBSS_BLOCK_LENGTH,BYTE] FIELD (BCBSZ_FIELDS)
%;
```

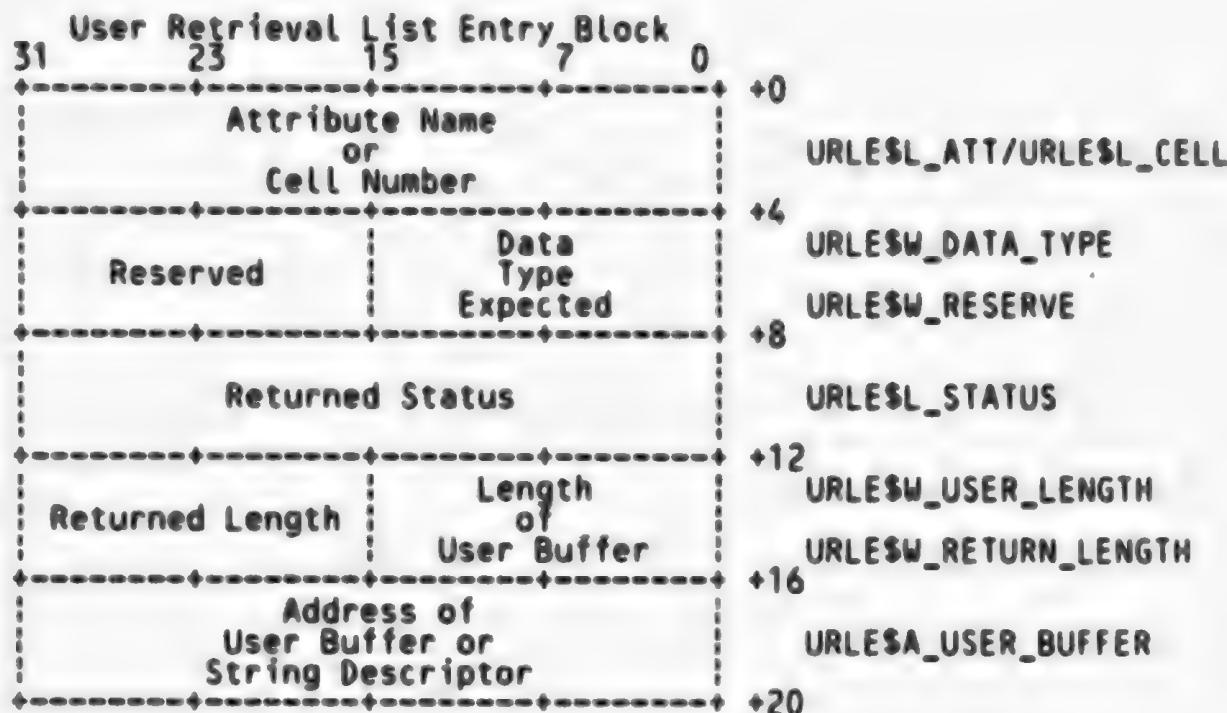
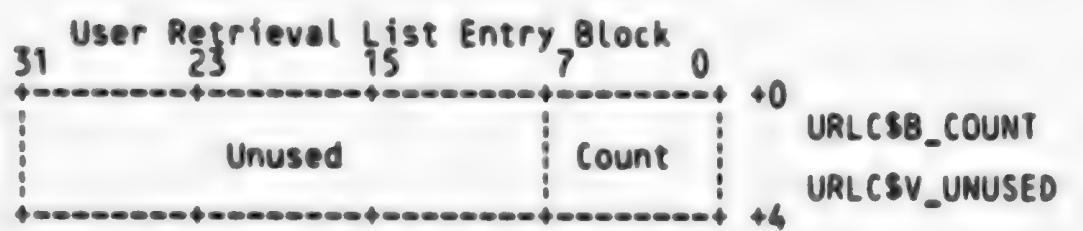
FIELD BCBSZ_FIELDS =

```
SET
BCBSA_NEXT      = [0, 0, 32, 0],
BCBSA_PRIOR     = [4, 0, 32, 0],
BCBSA_BUFFER    = [8, 0, 32, 0],
BCBSA_PCB       = [12, 0, 32, 0]
```

```
TES;
```

LITERAL

```
BCBSK_BUFFER_LIST = BLOCK[0, BCBSA_NEXT: , BYTE];
```



LITERAL

```
URLCSS_BLOCK_LENGTH = 4
URLESS_BLOCK_LENGTH = 20;
```

MACRO

```
SURLC = BLOCK[URLCSS_BLOCK_LENGTH,BYTE] FIELD (URLCSZ_FIELDS)
%;
```

FIELD URLCSZ_FIELDS =

```
SET URLCSB_COUNT = [0, 0, 8, 0]
URLCSV_UNUSED = [1, 0, 24, 0]
TES;
```

MACRO

```
SURLE = BLOCK[URLESS_BLOCK_LENGTH,BYTE] FIELD (URLESZ_FIELDS)
%;
```

FIELD URLESZ_FIELDS =

```
SET
```

```
URLESL_ATT      = [0, 0, 32, 0];
URLESL_CELL     = [0, 0, 32, 0];
URLESW_DATA_TYPE = [4, 0, 16, 0];
URLESW_RESERVE  = [6, 0, 16, 0];
URLESL_STATUS   = [8, 0, 32, 0];
URLESW_USER_LENGTH = [12, 0, 16, 0];
URLESW_RETURN_LENGTH = [14, 0, 16, 0];
URLESA_USER_BUFFER = [16, 0, 32, 0];
```

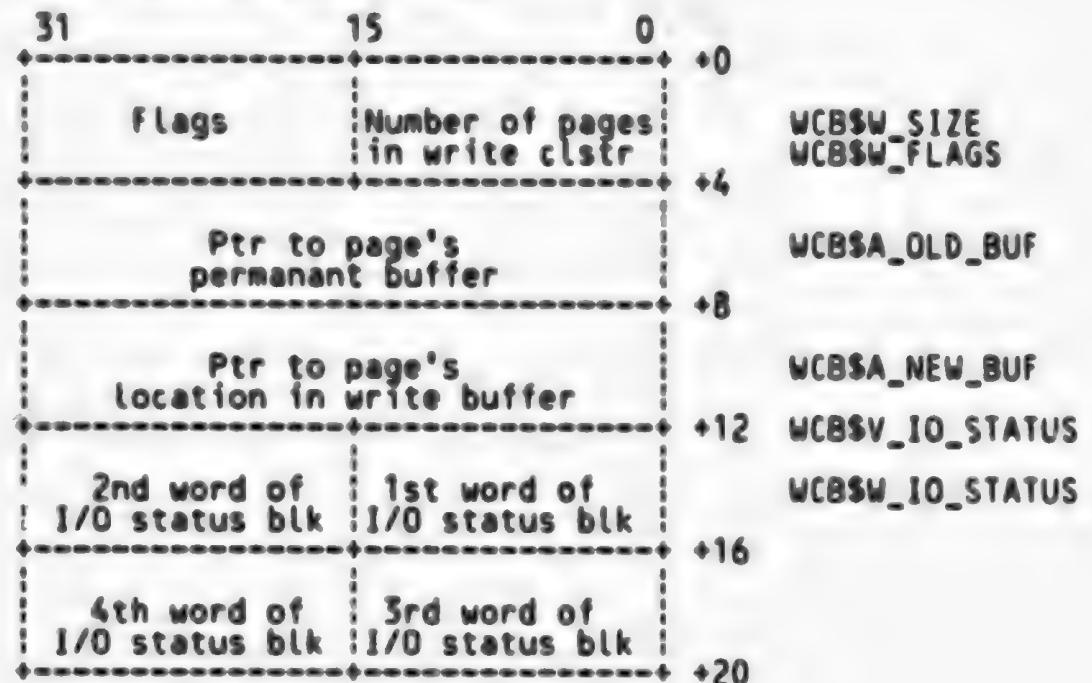
TES;

STRUCTURE

```
URLSBLOCK [I] = (URLSBLOCK + URLCSS_BLOCK_LENGTH +
                   (I * URLESS_BLOCK_LENGTH));
```

!+

Write Control Block (WCB)



!-

LITERAL
 WCBS_BLOCK_LENGTH = 20;

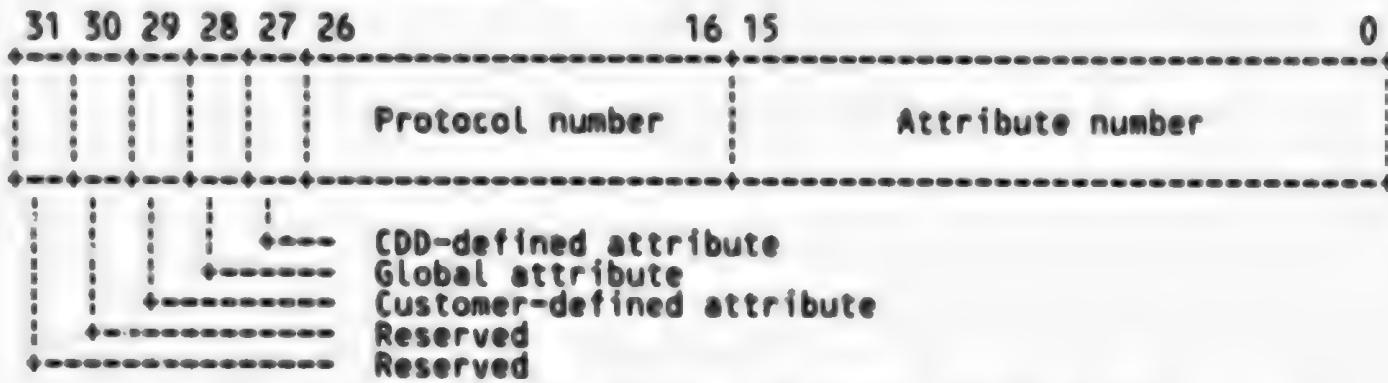
MACRO
 \$WCB = BLOCK[WCBS_BLOCK_LENGTH, BYTE] FIELD (WCBSZ_FIELDS)
 %;

FIELD SET WCBSZ_FIELDS =
 WCBS_SIZE = [0, 0, 16, 0],
 WCBS_FLAGS = [2, 0, 16, 0],
 WCBSV_WRITABLE = [2, 0, 1, 0], ! Write group to file
 WCBSV_GROUP = [2, 1, 1, 0], ! Group of pages
 WCBSA_OLD_BUF = [4, 0, 32, 0],
 WCBSA_NEW_BUF = [8, 0, 32, 0],
 WCBSV_IO_STATUS = [12, 0, 0, 0],
 WCBSW_IO_STATUS = [12, 0, 16, 0]

TES;

LITERAL
 WCBSM_WRITABLE = 1^1 - 1^0; ! Write group to file
 WCBSM_GROUP = 1^2 - 1^1; ! Group of pages

Attribute Name Block (ATNM)



The attribute name block defines the break down of the attribute name. Bits 0 to 15 contains the number. Bits 16 to 26 contains the protocol. The remaining bits are flags define the type of attribute. Bit 27 on implies a system defined attribute. Bit 28 on implies a global defined attribute. Bit 29 on implies a customer defined attribute.

NOTE: If the high order word of the attribute name block is 0 the block is describing a cell. That is the protocol will equal 0 and the flag bits will equal 0.

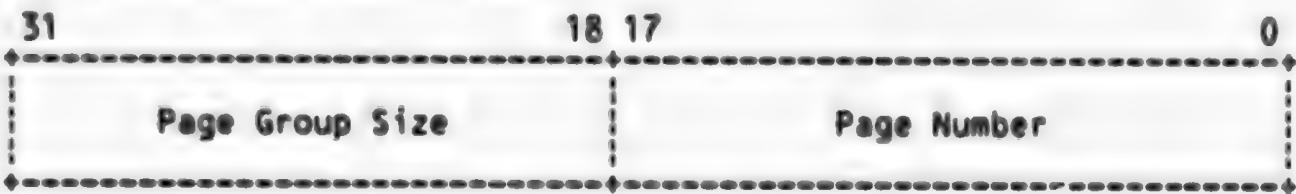
LITERAL
ATNMSS_BLOCK_LENGTH = 4;

MACRO SATNM = BLOCK[ATNMSS_BLOCK_LENGTH,BYTE] FIELD (ATNMSZ_FIELDS)
%;

FIELD SET ATNMSZ_FIELDS =
ATNMSW_NUMBER = [0, 0, 16, 0]
ATNMSW_HIGH_ORD = [0, 16, 16, 0]
ATNMSV_PROTOCOL = [0, 16, 11, 0]
ATNMSV_FLAGS = [0, 27, 5, 0]
ATNMSV_SYSTEM = [0, 27, 1, 0]
ATNMSV_GLOBAL = [0, 28, 1, 0]
ATNMSV_CUSTOMER = [0, 29, 1, 0]

YES:

* Page Number Block (PNB)



The page number block define the break down of the page number.
Bits 0 to 17 is the page number. Bits 18 to 31 define the number
of consecutive pages that can be read or written at one time.

LITERAL

```
PNBSS_BLOCK_LENGTH = 4;
```

MACRO

```
$PNB = BLOCK[PNBSS_BLOCK_LENGTH, BYTE] FIELD (PNBSZ_FIELDS)
```

```
%;
```

FIELD PNBSS_BLOCK_LENGTH =

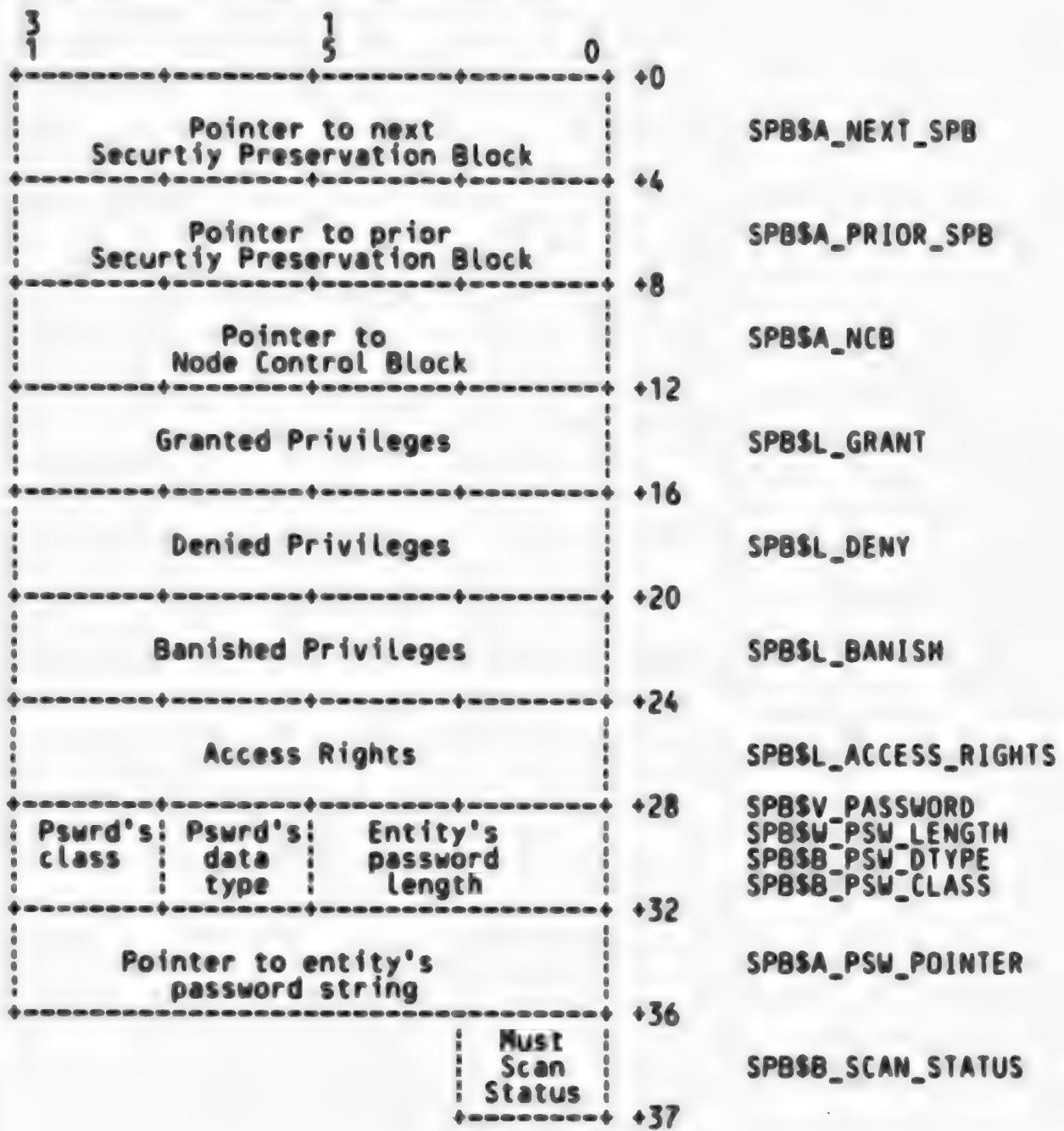
```
SET
```

```
PNBSV_PAGE_NUM = [0, 0, 17, 0]
```

```
PNBSV_GROUP_SIZ = [0, 17, 15, 0]
```

```
TES:
```

* Security Preservation Block (SPB)



LITERAL
SPB\$\$_BLOCK_LENGTH = 37;

MACRO

\$\$PB = BLOCK[SPB\$\$_BLOCK_LENGTH, BYTE] FIELD (SPB\$\$_FIELDS)

%:
FIELD SPB\$\$_FIELDS =

SET
SPBSA_NEXT_SPB = [0, 0, 32, 0],
SPBSA_PRIOR_SPB = [4, 0, 32, 0],
SPBSA_NCB = [8, 0, 32, 0],
SPBSL_GRANT = [12, 0, 32, 0],
SPBSL_DENY = [16, 0, 32, 0],
SPBSL_BANISH = [20, 0, 32, 0],
SPBSL_ACCESS_RIGHTS = [24, 0, 32, 0],
SPBSV_PASSWORD = [28, 0, 0, 0],
SPBSW_PSW_LENGTH = [28, 0, 16, 0],
SPBSB_PSW_DTYPE = [30, 0, 8, 0],
SPBSB_PSW_CLASS = [31, 0, 8, 0],
SPBSA_PSW_POINTER = [32, 0, 32, 0],
SPBSB_SCAN_STATUS = [36, 0, 8, 0]

TES;

LITERAL
SPBSK_QUE_HEADER = BLOCK[0, SPBSA_NEXT_SPB: , BYTE];

XSBTTL 'System Literal Definitions'

++ SYSTEM LITERAL DEFINITIONS

-- These literals are only used internally.

++ CDD Implementation Version

LITERAL

CDDSK_FACILITY	= 63,	
CDDSK_LOWEST VERSION	= 201,	! Lowest compatible version
CDDSK_VERSION	= 201;	! Present version

-- Boolean literals

LITERAL

TRUE	= 1;	! Boolean TRUE value
FALSE	= 0;	! Boolean FALSE value

-- The following literals are used to validate routines' parameter lists.

LITERAL

ARGSK_OPTIONAL	= 1,	! Parameter is optional
ARGSK_REQUIRED	= 2,	! Parameter is required
ARGSK_SYNC	= 3,	! See CDDSSU_VALIDATE documentation
ARGSK_SYNCIF	= 4,	! See CDDSSU_VALIDATE documentation
ARGSK_MARK	= 0,	! Mark parameter as missing
ARGSK_DEFAULT	= 1,	! Use default value
ARGSK_STRING	= 1,	! Default value is a null string
ARGSK_LONG	= 2,	! Default value is a longword
ARGSK_WORD	= 3,	! Default value is a word value
ARGSK_REF	= 0,	! Parameter passed by reference
ARGSK_VALUE	= 1;	! Parameter passed by value

-- Access Lock Constants

NOTE: The order of the LOCKSK_NORMAL lock constants MUST be the same as the order of the CIBSV_REF_COUNTS lock ref count fields in the CIB.

LITERAL

LOCKSK_NULL	= 0,	! Not locked
LOCKSK_NORMAL	= 1,	! Base of partially queued locks
LOCKSK_PRESENCE	= 1,	
LOCKSK_RETRIEVAL	= 2,	
LOCKSK_UPDATE	= 3,	
LOCKSK_DELETE	= 4,	
LOCKSK_NORMAL_TYPES	= 4,	! Number of partially queued locks
LOCKSK_QUEUED	= 5,	! Base of fully queued locks
LOCKSK_FULL_RET	= 5,	
LOCKSK_FULL_UPD	= 6,	
LOCKSK_QUEUED_TYPES	= 2,	! Number of fully queued locks
LOCKSK_XXX_GET	= 1^17-1^16,	! Flag indicates must establish lock
LOCKSK_XXX_GET_IF	= 1^18-1^17,	! Get lock if not already present
LOCKSK_OPTIONS	= 1^32-1^16:	! Option bits for locks

Types of page purge requests

LITERAL

PURGESK_ALL	= 1,	! Complete purge
PURGESK_PRESENCE	= 2,	! Purge but keep portal page
PURGESK_RETRIEVAL	= 3,	! Checkpoint & keep retrieval locks
PURGESK_UPDATE	= 4,	! Checkpoint & keep all locks
PURGESK_SUBTREE	= 100,	! Purge whole subtree
PURGESK_CLUSTER	= 101;	! Purge this cluster only

Types of blocks that can be allocated in a pool.

Types of pools

LITERAL

MEMSK_LOWEST_POOL	= 1,	! Lowest pool type
MEMSK_CCB_POOL	= 1,	! Pool for CCB and HCBs
MEMSK_CIB_POOL	= 2,	! Pool for cluster blocks
MEMSK_HIGHEST_POOL	= 2;	! Highest pool type

CCB Pool block types

LITERAL

MEMSK_CCB_LOWEST	= 1,	! Lowest block type in CCB pool
MEMSK_CCB_CCB	= 1,	! CCB block

MEMSK_CCB_LOW_SLOT	= 2,	! Lowest block type in free block list
MEMSK_CCB_HCB	= 2,	! HCB block
MEMSK_CCB_HIGHEST	= 2:	! Highest block type in CCB pool

!+ Cluster Pool block types

LITERAL

MEMSK_CIB_LOWEST	= 1,	! Lowest block type in cluster pool
MEMSK_CIB_CIB	= 1,	CIB block
MEMSK_CIB_LOW_SLOT	= 2,	Lowest block type in free block list
MEMSK_CIB_PCB	= 2,	PCB,
MEMSK_CIB_LCB	= 3,	LCB,
MEMSK_CIB_LCCB	= 4,	LCCB,
MEMSK_CIB_ECCB	= 5,	ECCB,
MEMSK_CIB_NCB	= 6,	NCB,
MEMSK_CIB_SPB	= 7,	SPB,
MEMSK_CIB_HIGHEST	= 7:	Highest block type in cluster pool

!+ These flags tell the deletion routine how it is to handle the following cases:

DELSK_FAST says that pointers do not have to be cleaned up, as the block they reside in is going to be deleted.

DELSK_PRESERVE indicates that a directory node is merely to be emptied, and that its cluster is not to be deleted.

DELSK_SUBDICTIONARY says that sub-files are to have their contents deleted.

LITERAL

DELSK_FAST	= 1^1 - 1^0,
DELSK_PRESERVE	= 1^2 - 1^1,
DELSK_SUBDICTIONARY	= 1^3 - 1^2:

!+ User Identification Criteria

LITERAL

CDDSK_ACL_LOWEST	= 1,	
CDDSK_ACL_PASSWORD	= 1,	! PASSWORD
CDDSK_ACL_TERMINAL	= 2,	! TERMINAL name or class
CDDSK_ACL_UIC	= 3,	! UIC
CDDSK_ACL_USERNAME	= 4,	! USERNAME
CDDSK_ACL_HIGHEST	= 4:	

%SBTTL 'Security Masks'

++
SECURITY MASKS

+- CDD security bits

LITERAL

CDDSK_PROT_C	= 1^1 - 1^0,	! CONTROL access
CDDSK_PROT_D	= 1^2 - 1^1,	! LOCAL DELETE access
CDDSK_PROT_G	= 1^3 - 1^2,	! GLOBAL DELETE access
CDDSK_PROT_H	= 1^4 - 1^3,	! HISTORY list entry creation access
CDDSK_PROT_P	= 1^5 - 1^4,	! PASS THRU access
CDDSK_PROT_S	= 1^6 - 1^5,	! SEE (read) access
CDDSK_PROT_U	= 1^7 - 1^6,	! UPDATE terminal node access
CDDSK_PROT_X	= 1^8 - 1^7,	! EXTEND directory node access
CDDSK_PROT_F	= 1^9 - 1^8,	! FORWARDing directory creation allowed

+- Macro-security values

CDDSK_PROT_ANY	= 1^9 - 1^0,
CDDSK_PROT_DELETE	= CDDSK_PROT_D OR CDDSK_PROT_G,
CDDSK_PROT_EXTEND	= CDDSK_PROT_F OR CDDSK_PROT_X,
CDDSK_PROT_UPDATE	= CDDSK_PROT_C OR CDDSK_PROT_D OR CDDSK_PROT_G OR CDDSK_PROT_H OR CDDSK_PROT_U OR CDDSK_PROT_X OR CDDSK_PROT_F,

+- Other processor security bits

+- VAX-11 Datatrieve

CDDSK_DTR_PROT_E	= 1^17 - 1^16,	! EXTEND file
CDDSK_DTR_PROT_R	= 1^18 - 1^17,	! READ file
CDDSK_DTR_PROT_M	= 1^19 - 1^18,	! MODIFY file
CDDSK_DTR_PROT_W	= 1^20 - 1^19:	! WRITE file

%SBTTL 'User Literal Definitions'

USER LITERAL DEFINITIONS

These symbols are needed by users of the program interface.

System Defined Attribute Names

LITERAL

CDDSK_SYSNAM_FLAGS = 1^28 OR 1^27 OR 0^16; ! Global/System-defined/Protocol=0

LITERAL

CDDSK_FIRST_SYSNAM	= 1 OR CDDSK_SYSNAM_FLAGS,	Lowest system defined attribute name value
CDDSK_FILE	= 1 OR CDDSK_SYSNAM_FLAGS,	Node's file name
CDDSK_HISTORY	= 2 OR CDDSK_SYSNAM_FLAGS,	History list head
CDDSK_NAME	= 3 OR CDDSK_SYSNAM_FLAGS,	Node's name
CDDSK_PROTOCOL	= 5 OR CDDSK_SYSNAM_FLAGS,	Node's protocol name
CDDSK_TYPE	= 6 OR CDDSK_SYSNAM_FLAGS,	Type of object pointed to by location code
CDDSK_PATHNAME	= 7 OR CDDSK_SYSNAM_FLAGS,	Node's complete pathname
CDDSK_SHORT_PATHNAME	= 8 OR CDDSK_SYSNAM_FLAGS,	Node's path to CDD\$DEFAULT directory
CDDSK_ORDER	= 9 OR CDDSK_SYSNAM_FLAGS,	Directory's order
CDDSK_LAST_SYSNAM	= 9 OR CDDSK_SYSNAM_FLAGS;	Highest system defined attribute name value

Attribute and Entity Types

LITERAL

CDDSK_FIRST_TYPE	= 1,
CDDSK_ENTITY	= 1,
CDDSK_ENTITY_LIST	= 2,
CDDSK_NULL	= 3,
CDDSK_NUMERIC	= 4,
CDDSK_STRING	= 5,
CDDSK_STRING_LIST	= 6,
CDDSK_DIRECTORY	= 7,
CDDSK_TERMINAL	= 8,
CDDSK_LAST_TYPE	= 8:

User's entity purge options

LITERAL

CDDSK_ALL	= 1^1 - 1^0,
CDDSK_ABORT	= 1^2 - 1^1,
CDDSK_CHECKPOINT	= 1^3 - 1^2;

!+ User's node creation options

LITERAL

CDDSK_NOHISTORY	= 1^1 - 1^0,	Doesn't want history list cluster
CDDSK_NOACL	= 1^2 - 1^1,	Don't create default ACL entry
CDDSK_CREATE	= 1^3 - 1^2,	Create dictionary file if needed
CDDSK_FIRST	= 1^4 - 1^3,	Insert as first node
CDDSK_LAST	= 1^5 - 1^4;	Insert as last node

!+ User's node deletion options

LITERAL

CDDSK_CHECK	= 1^1 - 1^0,	Fail if directory has children
CDDSK_SUBDICTIONARY	= 1^2 - 1^1;	Delete contents of subdictionaries

!+ Values of the CDDSK_ORDER attribute

LITERAL

CDDSK_SORTED	= 1,	Directory is sorted
CDDSK_NONSORTED	= 2;	Directory is not sorted

%SBTTL 'LINKAGE DEFINITIONS'

++
LINKAGE DEFINITIONS
--

|-+ CDDCALL

This linkage uses the CALLG/CALLS linkage convention, except that it allows for one global register to be used in parameter passing.

R11 -- used to pass the user's context pointer.

LINKAGE
CDDCALL = CALL : GLOBAL (USER_CONTEXT = 11);

|-+ SYS_JSB

This linkage provides us with a general JSB routine linkage.

LINKAGE
SYS_JSB = JSB;

```
** %SBTTL 'MACRO DEFINITIONS'
```

```
!-- MACRO DEFINITIONS
```

```
♦ $ACTIVE  
$INACTIVE
```

These macros declare that we have started, and finished, respectively, a CDD transaction. They abort the transaction if another transaction is in progress.

```
MACRO
```

```
  $ACTIVE =  
  BEGIN  
    EXTERNAL CDD$GB_INUSE:      BYTE;  
    EXTERNAL LITERAL CDD$_NOTASTREE;  
    BUILTIN TESTBITSS;  
  
    IF TESTBITSS (CDD$GB_INUSE) THEN  
      SIGNAL (CDD$_NOTASTREE);  
  END  
%,
```

```
  $INACTIVE =  
  BEGIN  
    EXTERNAL CDD$GB_INUSE:      BYTE;  
    CDD$GB_INUSE = FALSE;  
  END  
%,
```

```
♦ SBITLEAR
```

This macro checks to see if any bit in a mask is set in the target area. If not, it returns TRUE.

```
SBITLEAR(target, mask) =  
  (target AND mask) EQLU 0  
%,
```

\$BITSET

This macro checks to see if any bit in a mask is set in the target area. If so, it returns TRUE.

```
$BITSET(target, mask) =  
  (target AND mask) NEQU 0  
%.
```

\$DONE_TRANS

This macro is used to terminate a transaction.

Call:

```
$DONE_TRANS [(dsc1 [, dsc2] ...)]
```

Where:

dsc1 ::= the names of dynamic descriptors which are to have their strings returned to the string pool.

```
$DONE_TRANS (dsc1) =  
BEGIN  
  EXTERNAL ROUTINE  
    CDD$SN_DONE_TRANS      : CDDCALL      NOVALUE;  
  
    CDD$SN_DONE_TRANS (dsc1  
      %IF NOT %NULL(%REMAINING) %THEN , %REMAINING %FI );  
    .USER_CONTEXT[CCBSL_STATUS]  
  END  
%.
```

\$FIND_ENTITY

This macro returns the virtual address of the LCCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

Call:

```
lccb-block.wa.v = $FIND_ENTITY (valid-arg.ra.v ,  
  ( CHECK      { UPDATE      } ) , ( MODIFY      ) ;  
  RETRIEVAL      READ  
  DELETE        DELETE
```

ANY

Where:

valid-arg ::= the address of the calling routine's validated argument list.
 The argument list must have the following format:

valid-arg[0] ::= address of longword holding context #
 valid-arg[1] ::= address of descriptor holding path name, or zero.

The first set of keywords names the desired lock state of the entity's cluster.

The last set of keywords names the intended access to the cluster.

```

$FIND_ENTITY (valid_arg, locking, security) =
BEGIN
  EXTERNAL ROUTINE
    CDDSSN_FIND_ENTITY : CDDCALL;
    CDDSSN_FIND_ENTITY(..,valid_arg[1],
      $FIND_XXX_LOCKING (%REMOVE(locking)),
      %NAME('CDDSK_PROT_', security))
  END
%,

$FIND_XXX_LOCKING (class, type) =
%IF %IDENTICAL (class, %QUOTE LOCK) %THEN
  LOCKSK_XXX_GET OR
%ELSE
  %IF %IDENTICAL (class, %QUOTE LOCKIF) %THEN
    LOCKSK_XXX_GET_IF OR
  %ELSE
    %IF NOT %IDENTICAL (class, %QUOTE CHECK) %THEN
      %ERROR ('Invalid locking keyword: ', class)
    %FI
  %FI
%NAME ('LOCKSK_', type)
%,


```

\$FIND_NODE

This macro returns the virtual address of the NCB associated with a path name or location code. It also checks to make certain that the cluster is locked as requested, and that the target node allows the requested security access.

Call:

```
ncb-block.wa.v = SFIND NODE (valid-arg.ra.v ,
  LOCK          RETRIEVAL      READ
  ( ( LOCKIF ) ( UPDATE ) , ( MODIFY ) );
  CHECK         DELETE        DELETE
                ANY
```

Where:

valid-arg ::= the address of the calling routine's validated argument list.
The argument list must have the following format:

```
valid-arg[0] ::= address of longword holding context #
valid-arg[1] ::= address of descriptor holding path name, or zero.
valid-arg[2] ::= address of longword holding location code, or zero.
```

The first set of keywords names the desired lock state of the node's cluster.

The last set of keywords names the intended access to the cluster.

```
SFIND NODE (valid_arg, locking, security) =
BEGIN
  EXTERNAL ROUTINE
    CDDSSN_FIND_NODE      : CDDCALL;
  CDDSSN FIND NODE (valid_arg,
    SFIND XXX LOCKING (%REMOVE(locking)),
    %NAME? 'CDDSK_PROT_', security)
END
%.
```

SFIND_PARENT

This macro returns the virtual address of the NCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

Call:

```
status.wlc.v = SFIND PARENT (valid-arg.ra.v ,
  LOCK          RETRIEVAL      READ
  ( ( LOCKIF ) ( UPDATE ) , ( MODIFY ) , ncb-block.wa.r,
  CHECK         DELETE        DELETE
                ANY
  name.wt.ds);
```

Where:

valid-arg ::= the address of the calling routine's validated argument list.
The argument list must have the following format:

valid-arg[0] ::= address of longword holding context #
valid-arg[1] ::= address of descriptor holding path name, or zero.
valid-arg[2] ::= address of longword holding location code, or zero.

The first set of keywords names the desired lock state of the target cluster.

The last set of keywords names the intended access to the cluster.

```
SFIND_PARENT (valid_arg, locking, security, ncb_block, name) =
BEGIN
  EXTERNAL ROUTINE
    CDDSSSN_FIND_PARENT : CDDCALL;
  CDDSSSN_FIND_PARENT (valid_arg,
    SFIND_XXX_LOCKING (%REMOVE(locking)),
    %NAME%{CDD$K_PROT_%, security}, ncb_block, name)
END
.
```

SINIT_DSC

This macro is used to initialize dynamic string descriptors.

Call:

SINIT_DSC (dsc1 [, dsc2] ...)

```
SINIT_DSC[DSC_NAM] =
BEGIN
  DSC_NAM[DSC$B_DTYPE] = DSC$K_DTYPE_T;
  DSC_NAM[DSC$B_CLASS] = DSC$K_CLASS_D;
  DSC_NAM[DSC$W_LENGTH] = 0;
  DSC_NAM[DSC$A_POINTER] = 0;
END
.
```

SIO_SYNC (ef, iosb)

This macro waits for the event flag (ef) to be set and for the I/O status parameter (iosb) to be filled in (non-zero).

ef must be the target event flag number.

iosb must be the address of the I/O status block. This must be defined as a VECTOR[WORD] VOLATILE structure.

```
$.BIO_SYNC (ef, iosb) =  
  DO  
    BEGIN  
      LOCAL  
      STATUS:      LONG;  
  
      STATUS = SWAITFR (efn = ef);  
      IF NOT .STATUS THEN  
        SIGNAL STOP (.STATUS);  
      STATUS = $CLREF (efn = ef);  
      IF NOT .STATUS THEN  
        SIGNAL STOP (.STATUS);  
    END WHILE .iosb[0] EQLU 0
```

%,

SMARK_PAGE (page-block)

This macro marks a page as modified. Such a page must be written back to the dictionary file when it is purged from the staging buffers.

```
$.SMARK_PAGE (page_block) =  
  page_block[PCBSV_MODIFIED] = TRUE
```

SPARAMETERS (arg1 [, arg2] ...)

This macro is used to build the control vector for the parameter list validate routine (CDDSSU_VALIDATE).

There is one entry (arg1, arg2, etc) for every formal parameter in the routine. Each entry has the following format:

```
( REQUIRED [n] , [REF : VALUE] :  
  { SYNC : $YNCIF } , n :  
  OPTIONAL [,MARK : ,DEFAULT [,STRING : ,LONG : ,WORD] ] )
```

```
$.SPARAMETERS[] =  
  UPLIT ($LENGTH, SPARM_PROCESS ($REMOVE ($REMAINING)))
```

%,

SPARM PROCESS[arg] =
 BTTE (\$PARM_DECIDE (%REMOVE(arg)))

%,

SPARM_DECIDE(status)[] =
 %IF %IDENTICAL (status, %QUOTE REQUIRED) %THEN
 ARGSK_REQUIRED, 0,
 %IF %NULL (%REMAINING) %THEN
 ARGSK_REF, 0
 %ELSE
 \$PARM_REQUIRED (%REMAINING)
 %FI
%ELSE
 %IF %IDENTICAL (status, %QUOTE SYNC) %THEN
 ARGSK_SYNC, 0, 0, %REMAINING
 %ELSE
 %IF %IDENTICAL (status, %QUOTE SYNCIF) %THEN
 ARGSK_SYNCIF, 0, 0, %REMAINING
 %ELSE
 %IF %IDENTICAL (status, %QUOTE OPTIONAL) %THEN
 ARGSK_OPTIONAL, \$PARM_OPTIONAL (%REMAINING), 0
 %ELSE
 %ERROR ('Invalid parameter status: ', status)
 0, 0, 0
 %FI
 %FI
%FI
%,

SPARM REQUIRED(number, pass) =
 %IF %NULL (pass) %THEN
 0
 %ELSE
 %NAME ('ARGSK_', pass)
 %FI
%IF %NULL (number) %THEN
 0
%ELSE
 , number
%FI

%,

SPARM OPTIONAL(action)[] =
 %IF %IDENTICAL (action, %QUOTE MARK) %THEN
 ARGSK_MARK, 0
 %ELSE
 %IF %IDENTICAL (action, %QUOTE DEFAULT) %THEN
 ARGSK_DEFAULT, %NAME ('ARGSK_', %REMAINING)
 %ELSE
 %ERROR ('Invalid action keyword: ', action)
 0, 0
 %FI
 %FI

X.

SPRESENT

This macro checks to see if a parameter is present (non-zero) in a list.

SPRESENT (name) =
.name NEQA 0

X.

SRECOVERY (RESET : ENABLE : DISABLE)

This macro determines the ability of the exit handler to perform a purge of the cache if the task aborts.

SRECOVERY (DISABLE)

declares that the internal data structures or external disk structure is in an indeterminant state and cannot be recovered by the exit handler.

SRECOVERY (ENABLE)

declares that the data structure manipulation is completed.

SRECOVERY (RESET)

specifies that the data structures are in a recoverable state.

Note that these can be nested. Recovery is only possible if the recovery counter is zero (reset).

SRECOVERY (option) =
BEGIN
 EXTERNAL
 CDDSGW_RECOVERY: WORD;
 %IF %IDENTICAL (option, %QUOTE DISABLE) %THEN
 CDDSGW_RECOVERY = .CDDSGW_RECOVERY + 1;
 %ELSE
 %IF %IDENTICAL (option, %QUOTE ENABLE) %THEN
 CDDSGW_RECOVERY = .CDDSGW_RECOVERY - 1;
 %ELSE
 %IF %IDENTICAL (option, %QUOTE RESET) %THEN
 CDDSGW_RECOVERY = 0;
 %ELSE
 %ERROR ('Illegal recovery option: ', option)
 %FI
 %FI
 %FI

%, END XFI

SRELEASE_LOCK

This macro calls the CDDSSN_RELEASE routine to release one or more locks.

Call:

```
SRELEASE_LOCK ( { RETRIEVAL
                  UPDATE }, ncb-block1 ... );
                  DELETE
```

```
SRELEASE_LOCK (locking)[] =
BEGIN
  EXTERNAL ROUTINE
    CDDSSN_RELEASE      : CDDCALL      NOVALUE:
    CDDSSN_RELEASE (%NAME ('LOCKSK_', locking), %REMAINING)
END
%,
```

SSIGNAL_SEVERE (error)

This routine signals a severe error.

```
SSIGNAL_SEVERE (ERROR) =
  SIGNAL (ERROR OR STSSK_SEVERE)
%,
```

SSTATIC_DSC

This macro is used to initialize static string descriptors.

Call:

```
SSTATIC_DSC (dsc1 [, dsc2] ...)
```

Where:

```
dsc1 ::= name ! ( name , source )
```

Source is the name of another descriptor. The named descriptor is initialized to point to the same string as source.

```

$STATIC_DSC[dsci] =
  $STATIC_DSC_BUILD (%REMOVE (dsci))
%.

$STATIC_DSC_BUILD(name, source) =
BEGIN
  name[DSC$B_DTYPE] = DSC$K_DTYPE_T;
  name[DSC$B_CLASS] = DSC$K_CLASS_S;
  %IF %NULL(source) %THEN
    name[DSC$W_LENGTH] = 0;
    name[DSC$A_POINTER] = 0;
  %ELSE
    name[DSC$W_LENGTH] = .source[DSC$W_LENGTH];
    name[DSC$A_POINTER] = .source[DSC$A_POINTER];
  %FI
END
%.

```

STRING

These macros are used to build string descriptors for literal strings.

Call:

\$\$STRING ((name, string) ...)
\$\$STRING_INIT ();

"name" is defined to be a **BLOCK** structure, and the address of the string is poked into the structure by the **STRING_INIT** macro, which must be the first executable statement in a routine.

%,
\$STRING_PTR_SETUP[PAIR] =
\$STRING_PTR_INIT (%REMOVE(PAIR))
%,
\$STRING_PTR_INIT (STR_NAME, STR_VAL) =
STR_NAME[DSC\$A_POINTER] = UPLIT BYTE(%REMOVE(STR_VAL));
%,

+
\$TEXTC

This macro is used to define counted strings. The first byte of such strings is a count of the number of characters in the string.

Each string is defined to be a VECTOR[BYTE] structure, with the 0 element being the character count, and the actual string starting at STRING[1].

\$TEXTC ((name1,'str1') [, (name2, 'str2')] ...);

-
\$TEXTC[PAIR] =
\$TEXTC_STR(%REMOVE(PAIR))
%,
\$TEXTC STR(NAME,TSTR) =
BIND
NAME = UPLIT BYTE (%CHARCOUNT(%REMOVE(TSTR)), %REMOVE(TSTR)) :
VECTOR[%CHARCOUNT(%REMOVE(TSTR))+1, BYTE];
%,

+
\$VALIDATE (cntl, [arg-list])

This macro generates a call to the general transaction setup routine.

cntl -- the address of the control vector for CDD\$U_VALIDATE.

arg-list-- is optional. If present, it is the address of the vector which is to receive the verified argument list from CDD\$U_VALIDATE.

-
\$VALIDATE(cntl, arg_list) =
BEGIN
EXTERNAL ROUTINE
CDD\$SN_START_TRANS : CDDCALL NOVALUE;

```
BUILTIN  
  AP;  
  
  %IF %NULL (arg_list) %THEN  
    CDDSSN_START_TRANS (.AP, cntl)  
  %ELSE  
    CDDSSN_START_TRANS (.AP, cntl, arg_list)  
  %FI  
END  
%;
```

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